

A formal analysis of Bluetooth device discovery.

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Outline

- Motivation
- Bluetooth Protocol
- Probabilistic model checking and PRISM
- Modelling
- Results

Motivation

- Bluetooth
 - increasingly common wireless communication protocol
 - performance is particularly important in low-power setting
 - performance is the result of a non-trivial interaction between devices
 - ➔ formal verification desirable
 - protocol is randomised
 - ➔ need probabilistic formal verification

Bluetooth protocol

Bluetooth overview

- short-range low-power wireless protocol
- frequency hopping over 79/32 frequencies

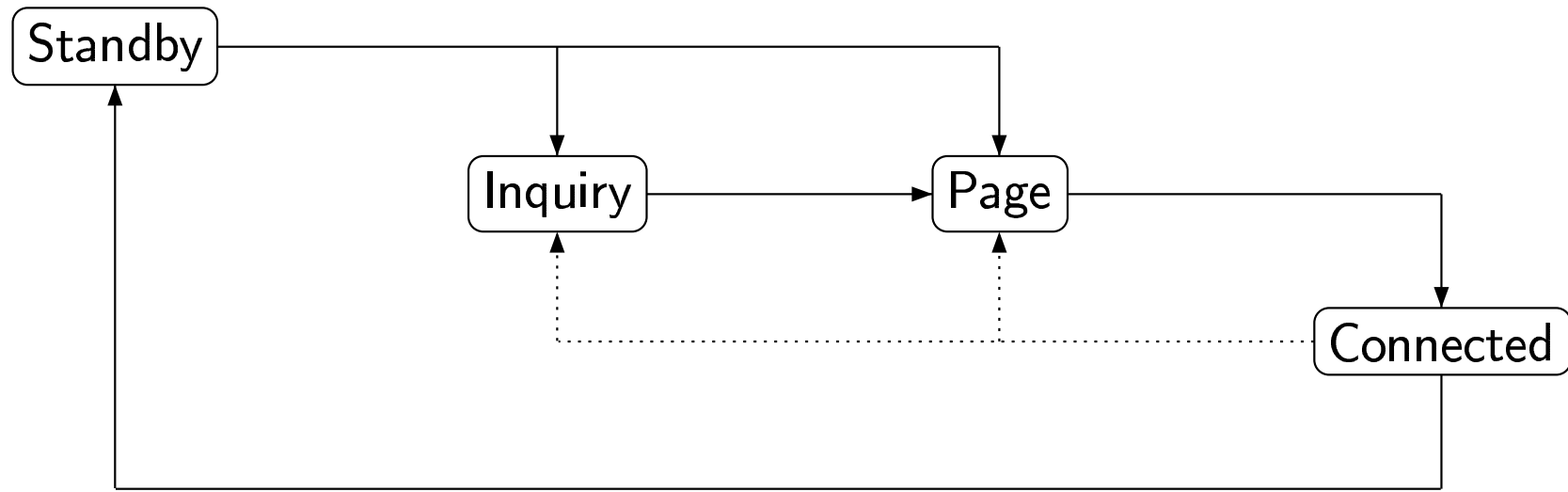
Bluetooth overview

- short-range low-power wireless protocol
- frequency hopping over 79/32 frequencies
 - need to form piconets
 - processes know when to send/receive
 - processes know the hopping frequency
 - master-slave roles
 - no communication before initialisation

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- First mandatory step: device discovery

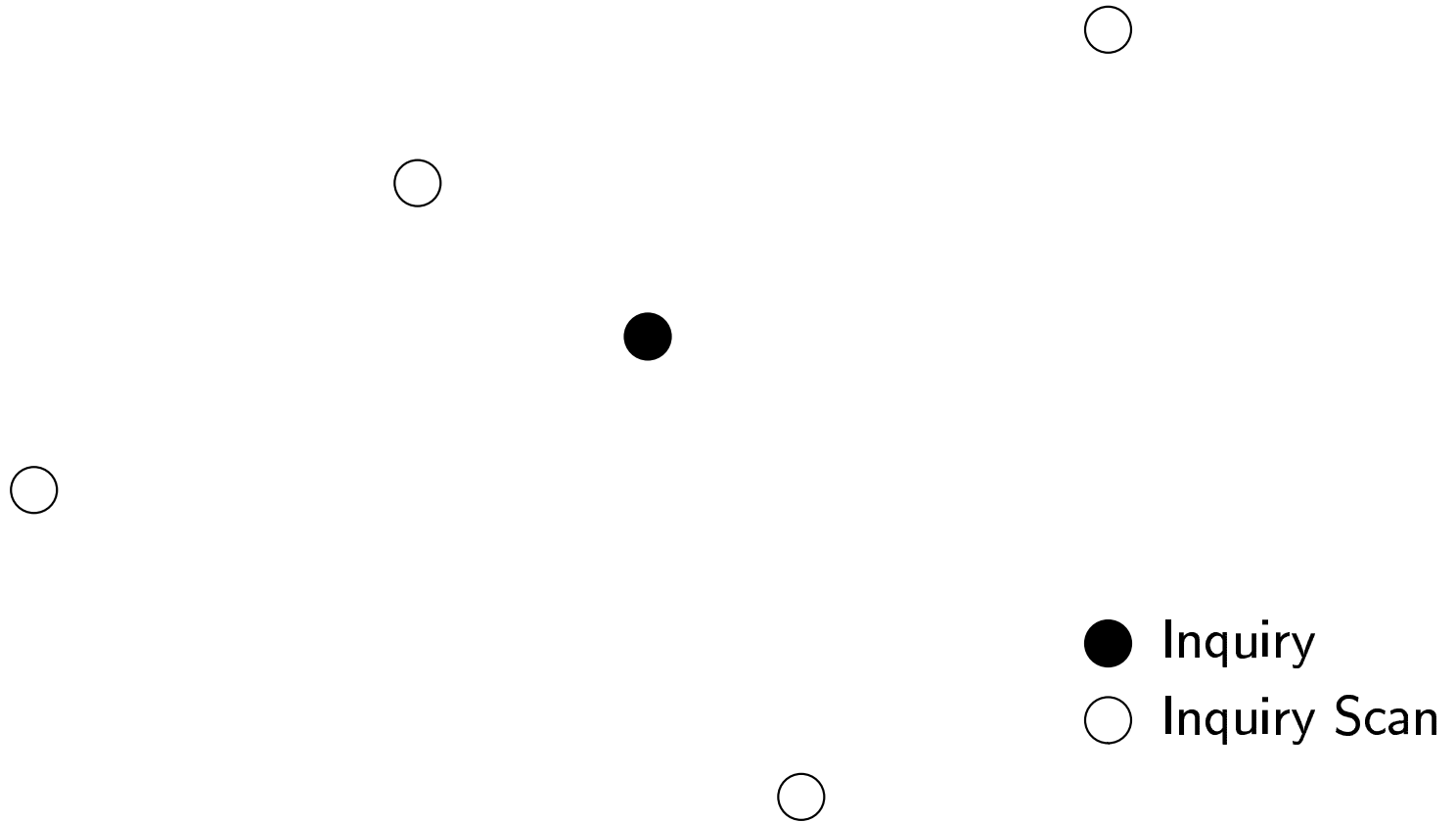
States of a Bluetooth device



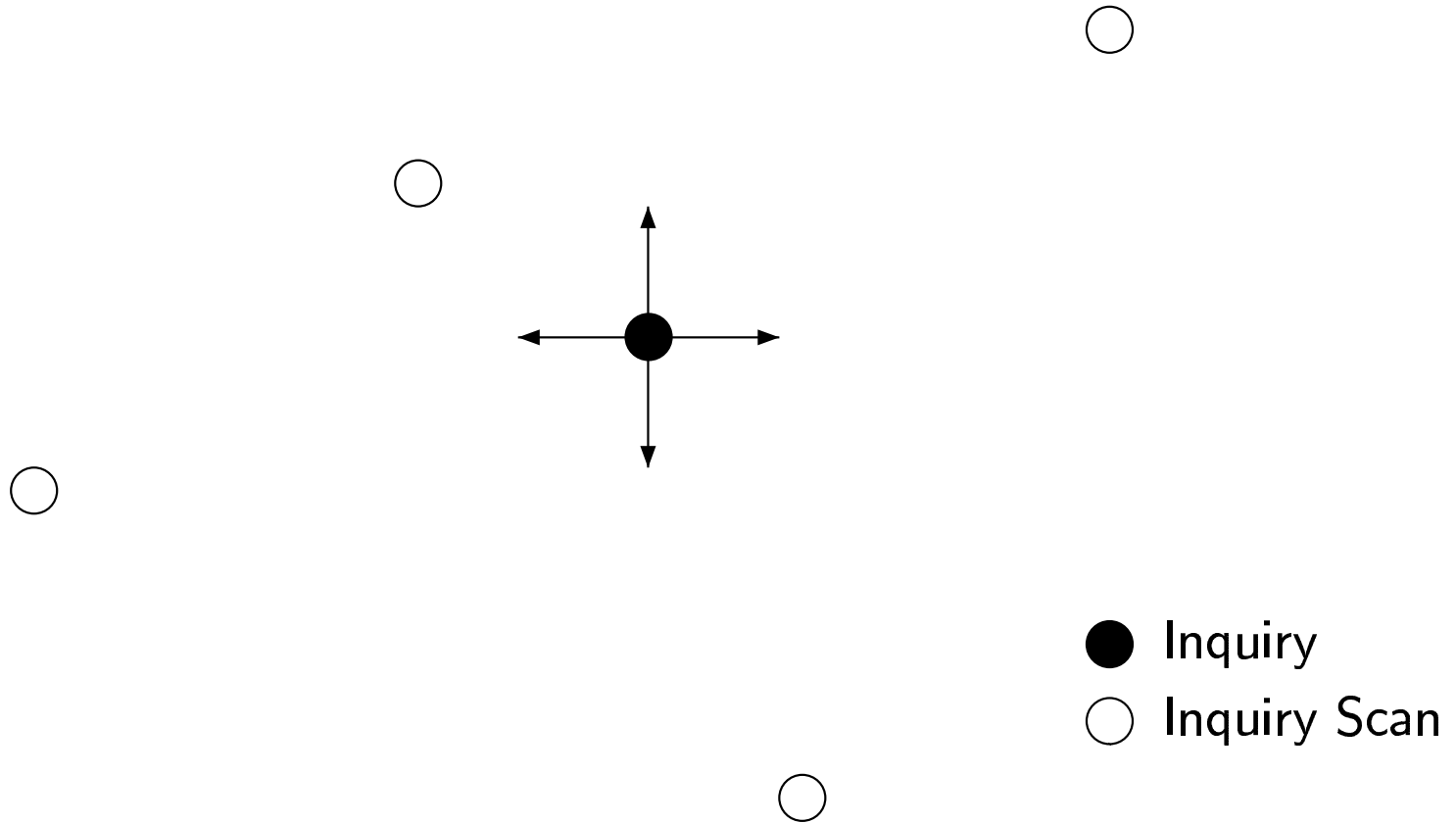
Standby: default operational state

Connected: device ready to communicate in a piconet

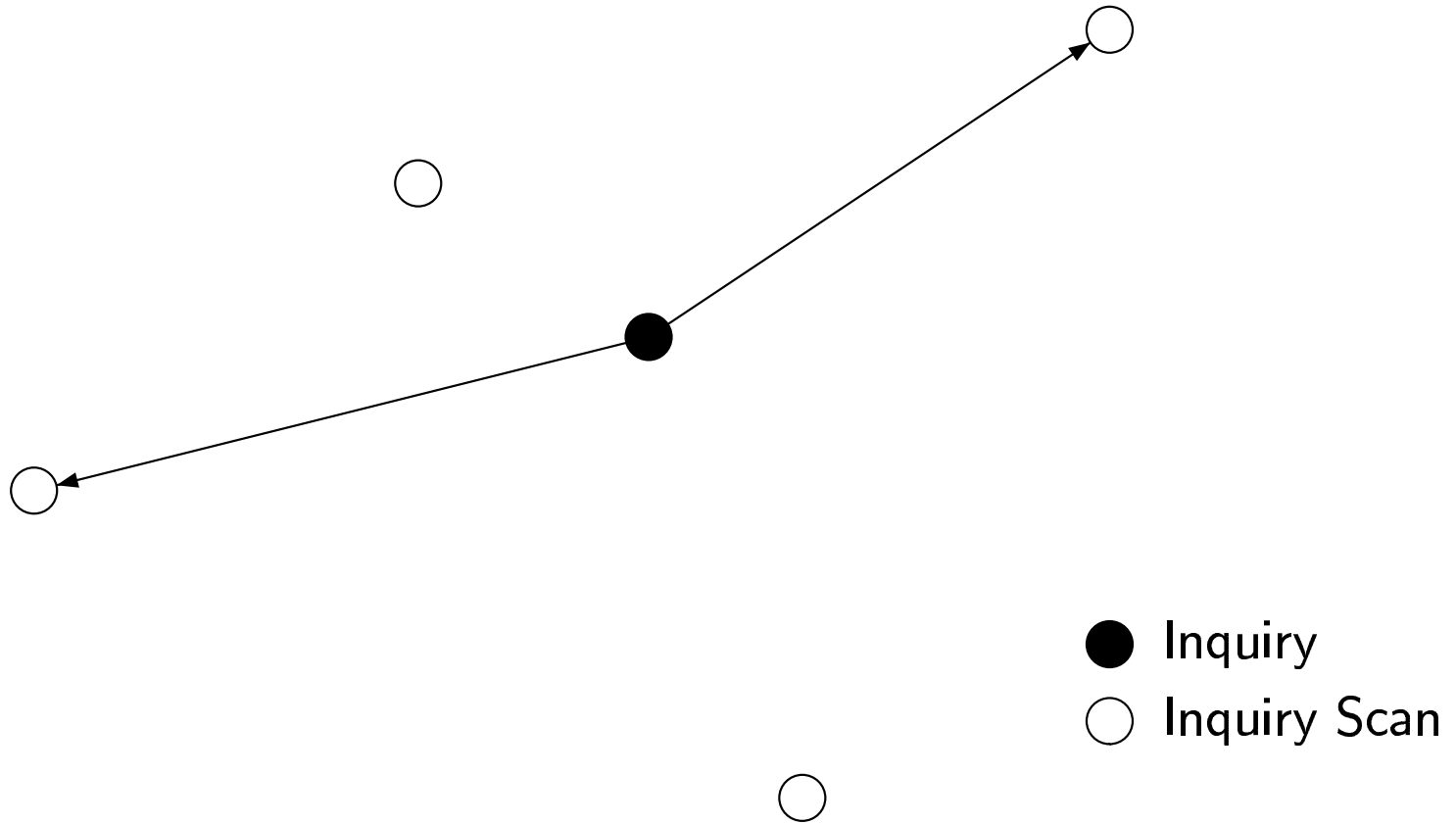
Inquiry (version 1.2)



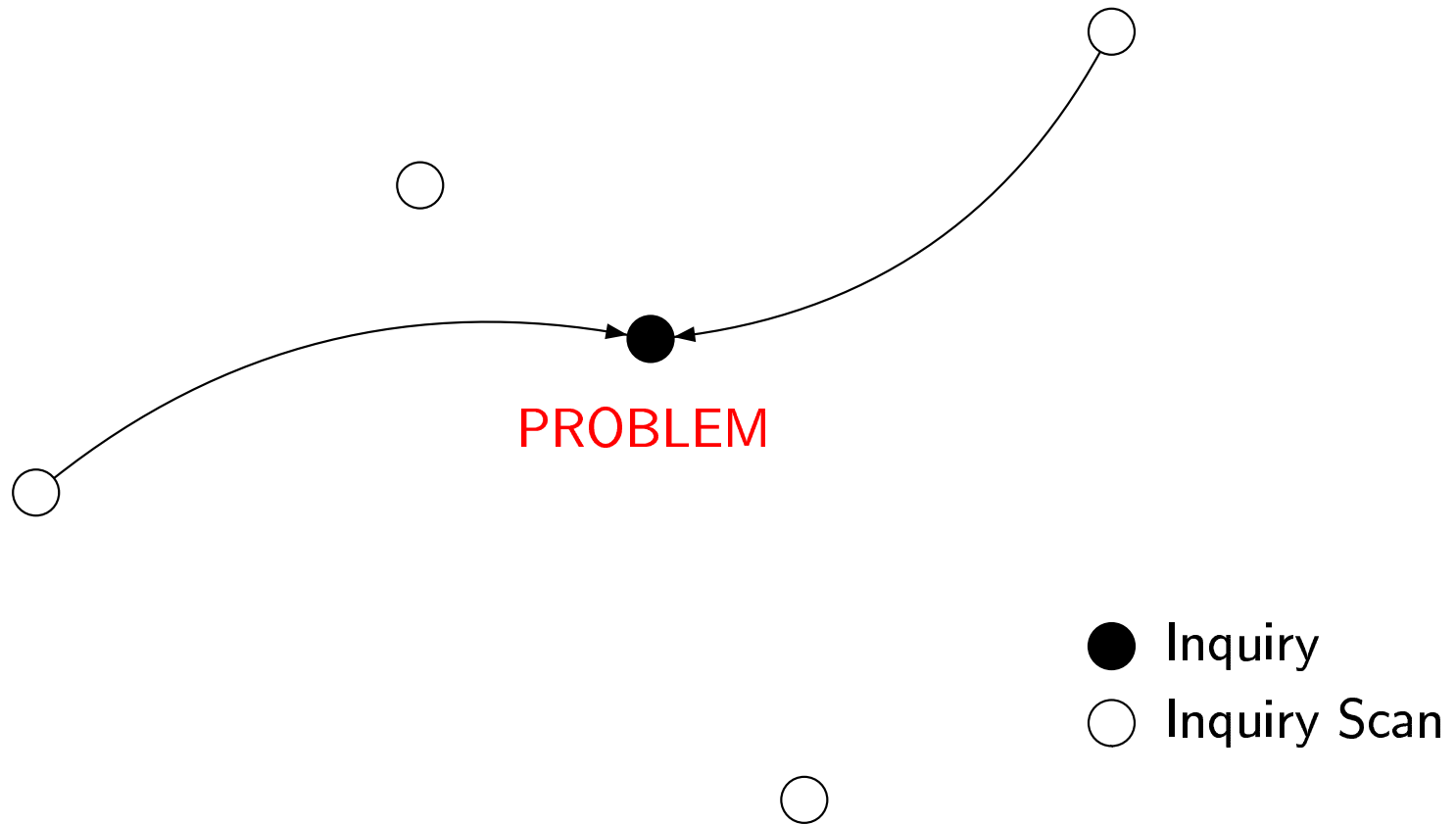
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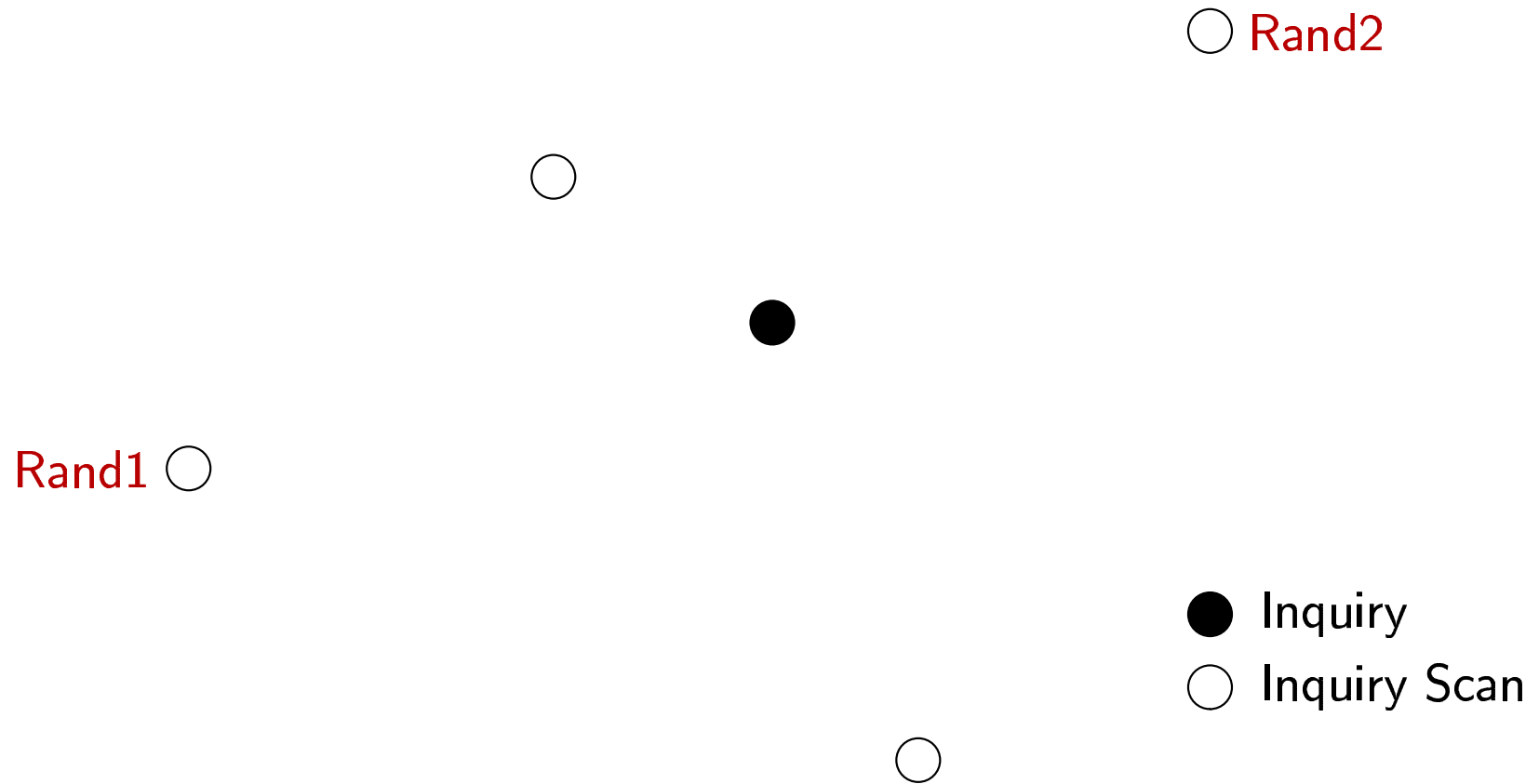
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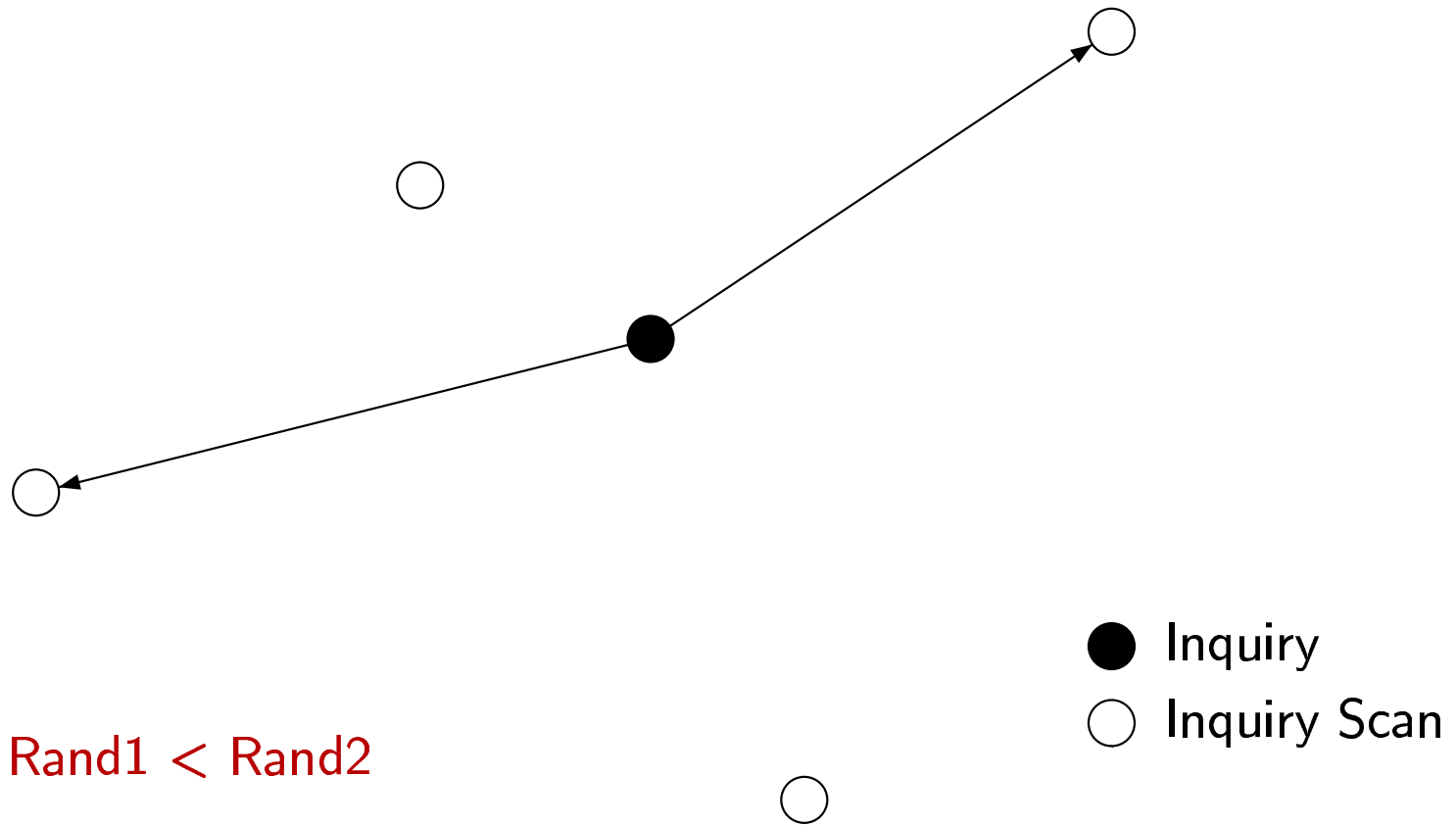
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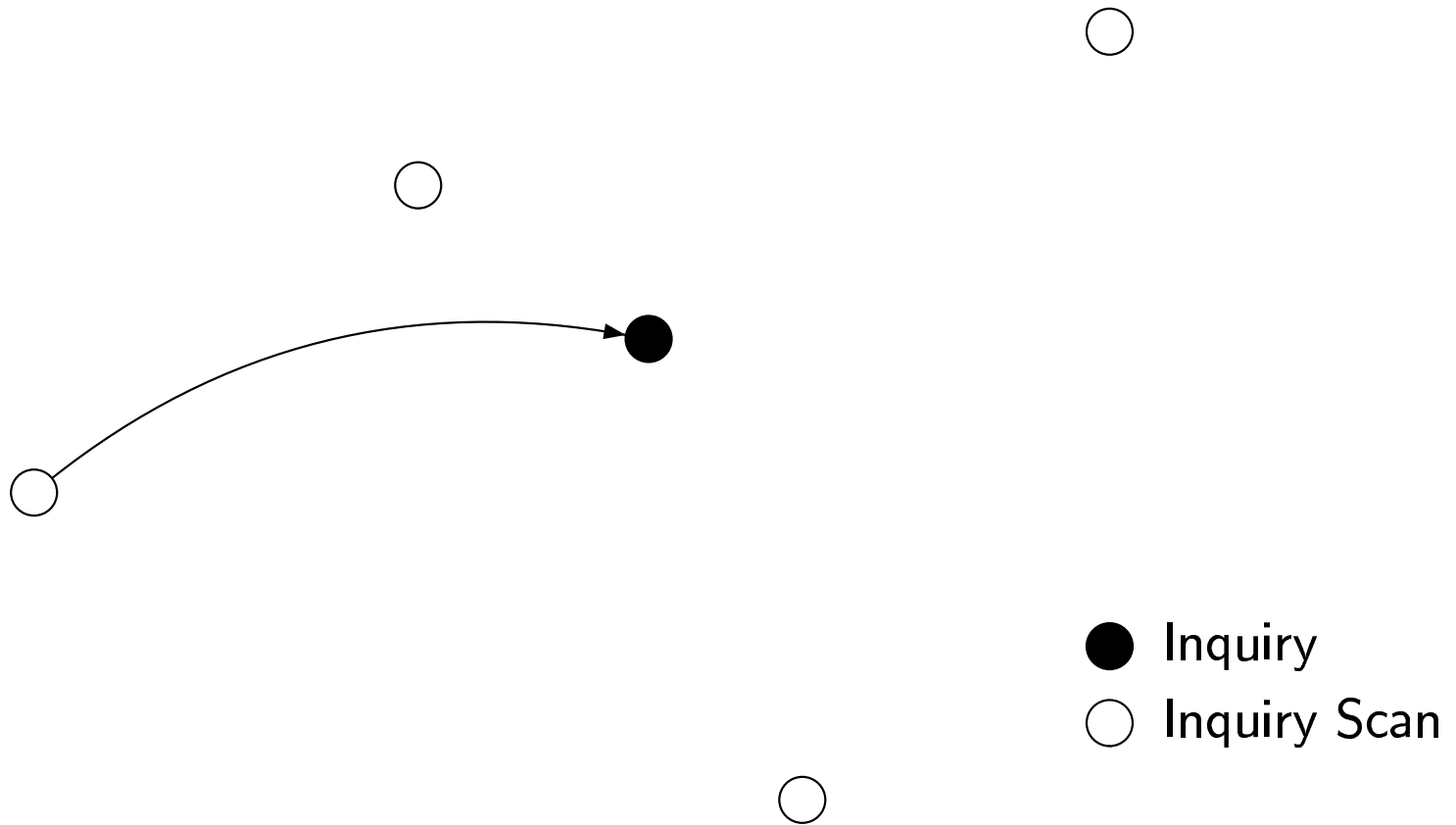
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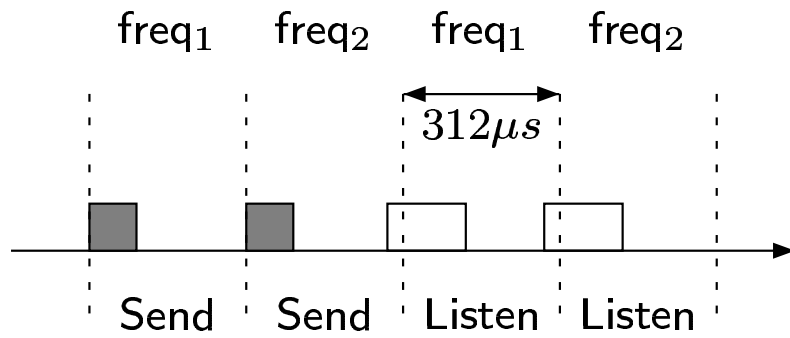
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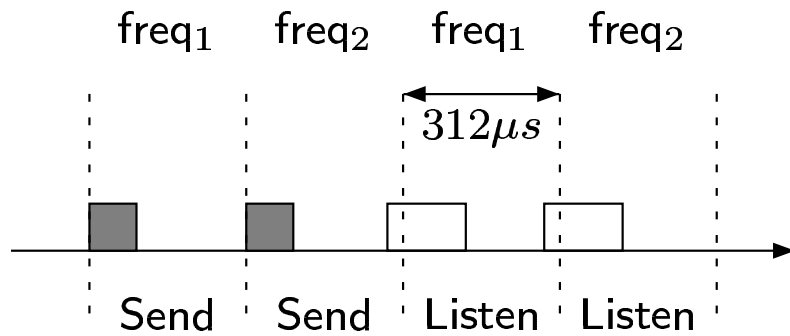
Inquiry (version 1.2)



The sender

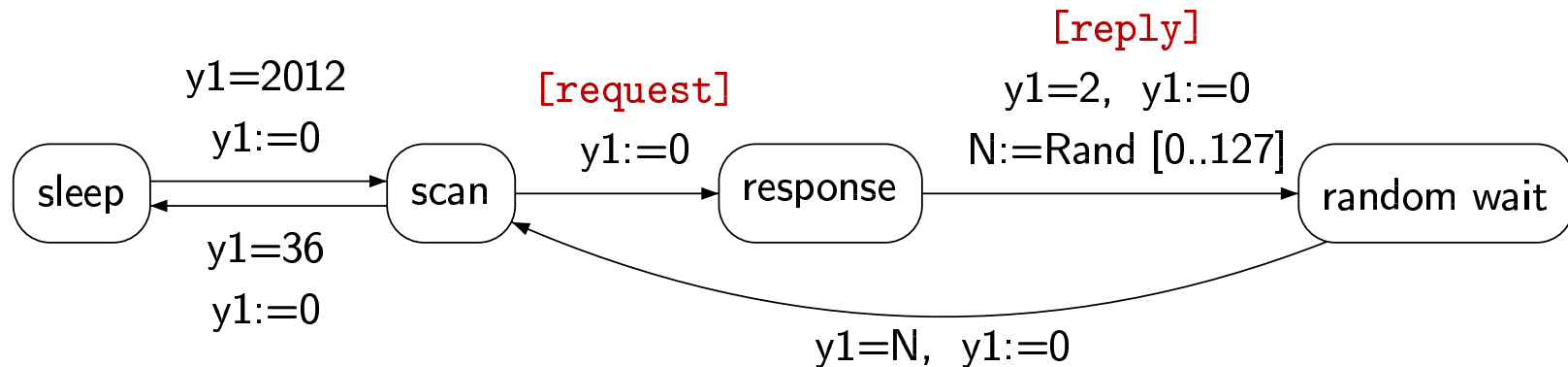


The sender



$$\text{freq} = [CLK_{16-12} + k \\ + (CLK_{4-2,0} - CLK_{16-12}) \bmod 16] \bmod 32$$

The receiver



- `[request]`: message sent by the sender
- `[reply]`: message sent by the receiver
- need to compute the frequency at which the receiver is listening (phase)
- phase increased by one each time the receiver replies

Probabilistic model checking and PRISM

Probabilistic model checking

- automatic formal verification of systems which exhibit stochastic behaviour
- conventional model checking:
 - model + property \longrightarrow yes/no
 - graph analysis
- **probabilistic** model checking:
 - **probabilistic** model + **probabilistic** property \longrightarrow yes/no/**probability**
 - graph analysis, numerical computation

Probabilistic models

- discrete-time Markov chains (DTMCs)
 - discrete time/probabilities
- continuous-time Markov chains (CTMCs)
 - real time (exponential distributions)
- Markov decision processes (MDPs)
 - discrete time/probabilities + nondeterminism

Probabilistic specifications

- PCTL/CSL - probabilistic extensions of CTL
- CTL: universal/existential quantification over paths
 - $AF\ success, EF\ success$
- PCTL: probabilistic quantification over paths
 - $P_{\geq 0.98}(F\ success)$
- Extension: Query actual probability values
 - $P_{=?}(F\ success)$
- Extension: Add rewards (or costs) to states/transitions of model...
 - $R_{=?}(F\ success)$
 - expected reward cumulated to reach *success*

PRISM

- PRISM: PRobabilistic Symbolic Model checker
- support for DTMCs/CTMCs/MDPs and PCTL/CSL
- high-level modelling language
 - based on Reactive Modules [Alur/Henzinger]
- wide range of existing case studies: randomised distributed algorithms, probabilistic security protocols, probabilistic communication protocols, etc.
- freely available: www.cs.bham.ac.uk/~dxp/prism

Modelling the protocol

Modelling formalism

- randomised back-off
 - need probabilistic model
 - discrete time slots (negligible clock drift)
 - no nondeterminism
 - no nondeterministic choice within a device
 - full synchronisation between devices
- ➔ discrete-time Markov Chains (DTMCs)

PRISM code: receiver's frequency

```
module frequency1

  z1 : [1..phase]; // clock for phase
  f1 : [1..16]; // frequency of receiver
  o1 : [0..1]; // offset of receiver

  // update frequency (1 slot passes)
  [time] z1 < phase -> (z1' = z1 + 1);
  [time] z1 = phase -> (z1' = 1) & (f1' = f1 < 16 ? f1 + 1 : 1) & (o1' = f1 < 16 ? o1 : 1 - o1);
  // update frequency: something is sent by the receiver
  [reply] true -> (f1' = (f1 < 16) ? f1 + 1 : 1) & (o1' = (f1 < 16) ? o1 : 1 - o1);

endmodule
```

State space explosion

- Sender changes state every time slot
- Receiver can wait for 2012 time slots without changing state
- 2 trains of 16 frequencies
- The trains change with time
- A train is repeated 256 times before switching
- The phase changes every 4096 slots

➤ A Huge model

➤ Too many possible initial states

➔ Need to use abstractions

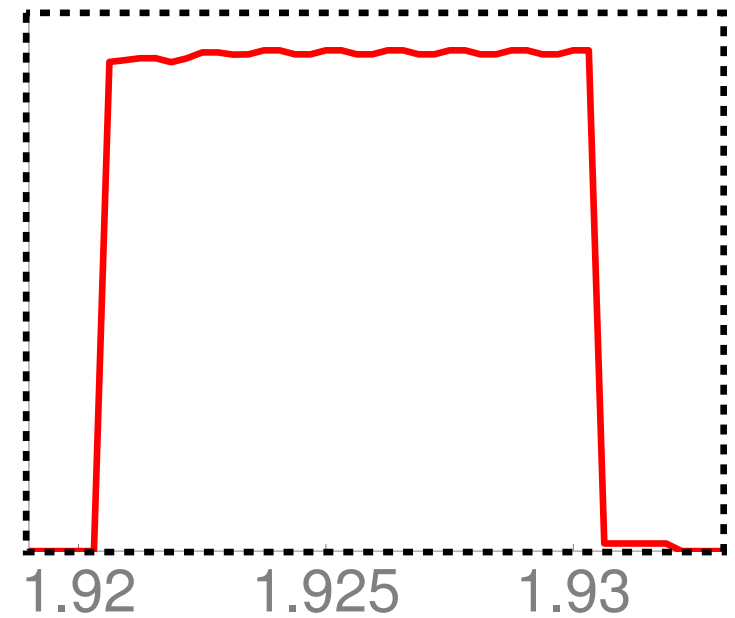
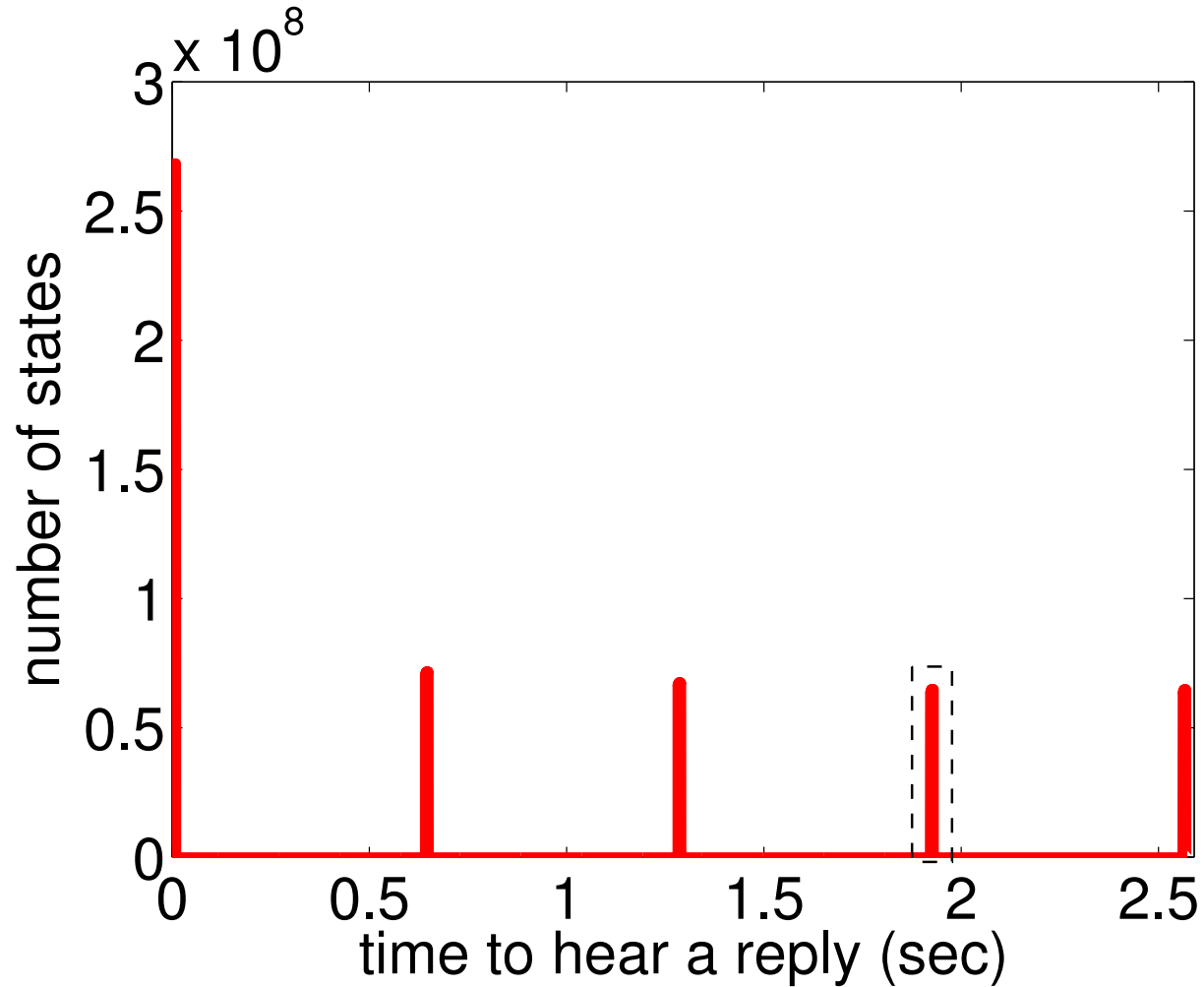
Verification and results

Expected time for one reply

- Big models → symbolic implementation (MTBDDs)
- Initial states split into 32 classes (possible initial frequencies)
- 32 models of around 3 thousand million states each
- 55-57 seconds to build one model
- 1-2 seconds to check the property

N	0	1	2	3	4
p	0.5003	0.6335	0.7591	0.8797	1

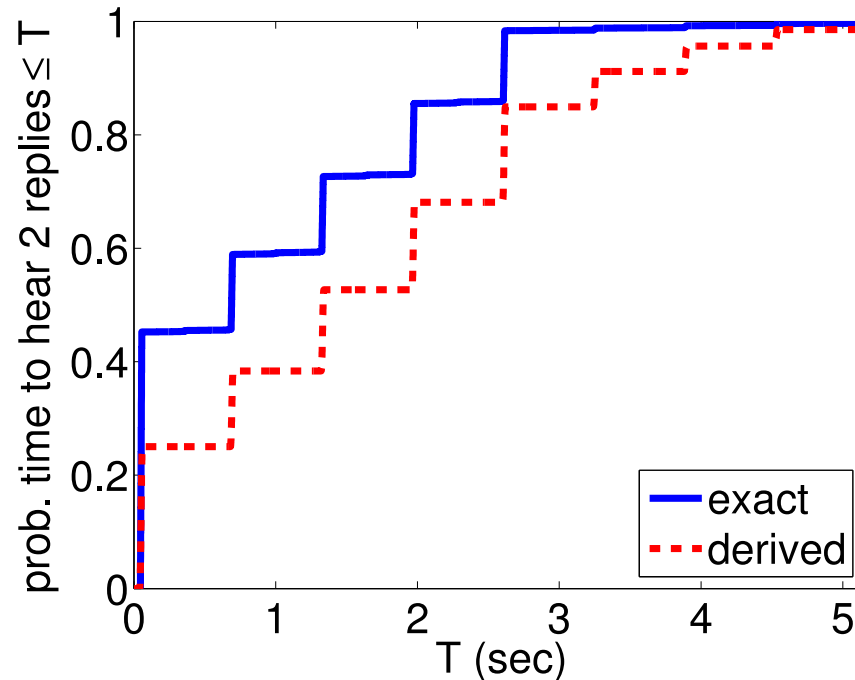
The graph for one reply



Expected time for two replies

- 32 models of 56 thousand million states each
- Time to build a model: 80 mins
- Time to verify the property : 165 mins
- Maximum expected time: 16565 slots (5 seconds)

The graph for two replies



- Need to really build the model for two replies
- Approximation via convolution is incorrect

Comparison with version 1.1

- In version 1.1, the receiver waits before sending a reply
- He replies to every second message.
- Expected time to receive a reply is longer

K	$n = 1$	$n = 1$ (v.1.1)	$n = 2$	$n = 2$ (derived)
0	0.500305	0.461240	0.455379	0.250305
1	0.633575	0.596265	0.590829	0.383657
2	0.759062	0.731585	0.728684	0.526981
3	0.879674	0.857913	0.855329	0.681114
4	1	0.984295	0.984218	0.849408
5	1	0.988269	0.988294	0.911750
6	1	0.992398	0.992514	0.956496
7	1	0.996294	0.996519	0.985521
8	1	1	1	1

Conclusions

- Summary
 - A formal analysis of Bluetooth device discovery
 - Quantitative quality of service results on a non simplified model
 - Model-checking vs simulation
- Future work
 - Further abstractions and symmetry.
 - Combine model checking and simulation.