

PRISM – An Overview



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PRISM – Overview

- **Probabilistic Symbolic Model Checker**
 - formal analysis of probabilistic systems
- **Developed:**
 - at the **University of Birmingham**
 - for the past 4½ years
 - by: **M. Kwiatkowska, G. Norman, D. Parker**
 - also: J. Meyer-Kayser, S. Gilmore, A. Hinton,
R. Downing

PRISM – Supported Models

- **Discrete-time Markov chains (DTMCs)**
 - discrete state transition probabilities
- **Continuous-time Markov chains (CTMCs)**
 - time modelled as exponential distributions
- **Markov decision processes (MDPs)**
 - nondeterminism + probabilities
 - e.g. distributed probabilistic systems

PRISM – Specification Formalisms

- Extensions of temporal logic **CTL**
- **PCTL** (for properties of DTMCs, MDPs)
 - “**Prob.** of leader eventually being elected is **1**”
 - “**Prob.** of error occurring within **k** steps is **< 0.01**”
- **CSL** (for properties of CTMCs)
 - “**Prob.** of queue being full within **1 hour** is **< 0.2**”
 - “Long-run **prob.** of server being down is **< 0.05**”

PRISM – Basic Functionality

- **Parse** model description (PRISM language)
- **Construct** probabilistic model
- Compute **reachable/deadlock** states

- **Parse** properties (temporal logic)
- **Model check** each property
 - Return **Yes/No** + actual **probability**

Recent Extensions

- **Process algebra operators** added
 - More flexible parallel composition
 - Translation from stochastic process algebra **PEPA**
- **Support for rewards/costs** added
 - State-based rewards/costs in probabilistic models
 - Verification of properties such as
 - “Expected number of steps before termination is...”
 - “Expected time to elect a leader is...”

The PRISM Language

- State-based model description language
 - based on Reactive Modules [Alur, Henzinger]
- Basic ingredients:
 - **modules** with local, integer-valued **variables**
 - defined by **guarded commands**
 - combine through **parallel composition**
- Also:
 - **synchronisation** over action labels (CSP-style)
 - **global variables**

Example

- Randomised self-stabilisation [Israeli-Jalfon]
- Ring of identical processes, each has a token
- Stable state - only one process has a token
- Compute:
 - minimum probability of stabilising
 - maximum expected time to stabilise

```

// Israeli-Jalfon algorithm
// dxp/gxn 10/06/02
//ring of size 3

nondeterministic

global q1 : [0..1] init 1;
global q2 : [0..1] init 1;
global q3 : [0..1] init 1;

module start
    start : [0..1] init 0;
    // end of initialization phase
    [] start=0 -> start'=1;

endmodule

module process1
    // initialization (non-deterministic choice as to value of bits)
    [] start=0 -> q1'=0;
    [] start=0 -> q1'=1;
    // the protocol
    [] start=1 & (q1=1) -> 0.5 : q1'=0 & q3'=1 + 0.5 : q1'=0 & q2'=1;

endmodule

// add further processes through renaming
module process2=process1[q1=q2, q2=q3, q3=q1] endmodule
module process3=process1[q1=q3, q2=q1, q3=q2] endmodule

```

PEPA Extension

- Module parallel composition alternatives:
 - $P_1 \parallel P_2$
 - synchronise over alphabets
 - $P_1 \parallel [a,b,c] P_2$
 - synchronise over a named set of actions
 - $P_1 \parallel\!\!\parallel P_2$
 - fully asynchronous
- Action hiding: $P \setminus \{a,b,c\}$

Case Studies

- Randomised distributed algorithms/protocols
 - Leader election algorithms
 - Consensus protocol
 - Byzantine agreement protocol
 - IEEE 1394 FireWire Root Contention
 - Crowds (anonymity) protocol
 - IEEE 802.11 Wireless LAN
- CTMC models
 - queues, networks, manufacturing systems

Implementation

- **Java** – User interfaces, high-level code
- **C++** – underlying data structures/engines
- **CUDD** – BDD/MTBDD library
- **JavaCC** – language parsers
- **Source code** – GPL, web site download
- **Platforms** – Solaris, Linux, ...

Future Plans

- Simulator
- GUI development
- Parallel/distributed implementations
- Abstraction, compositionality, ...