Formal and Efficient Synthesis for Continuous-Time Linear Stochastic Hybrid Processes

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Abstract—Stochastic processes are expressive mathematical tools for modeling real-world systems that are subject to uncertainty. It is hence crucial to be able to formally analyze the behavior of these processes, especially in safety-critical applications. Most of the existing formal methods are not designed for continuous-time processes, and those that are typically suffer from state explosion in practice. This work introduces a theoretical framework and a scalable computational method for formal analysis and control synthesis for switched diffusions, a class of stochastic models with linear dynamics that are continuous in both time and space domains; the focus is on safety with possible extensions to other properties. The proposed framework first constructs a finite abstraction in the form of an uncertain Markov process through discretization of both time and space domains. The errors caused by the discretization in each domain are formally characterized and cast into the abstraction model. Then, a strategy that maximizes the probability of the safety property and is robust against the errors is synthesized over the abstraction model. Finally, this robust strategy is mapped to a switching strategy for the stochastic processes that guarantees the safety property. The framework is demonstrated in three case studies, including one that illustrates the trade-off of the error contribution by the time and space discretization parameters.

Index Terms—Formal verification, formal synthesis, probabilistic model checking, stochastic hybrid systems, switched diffusions, formal abstractions, temporal logics.

I. INTRODUCTION

SWITCHED stochastic processes are powerful models for mathematical representation of real-world systems. These models, through switching between several Stochastic Differential Equations (SDEs), enable the inclusion of uncertainty as well as control, both of which are intrinsic in the physical laws of systems and in their interactions with the world [2]. Therefore, they are frequently used for modeling and analysis of systems in a wide range of domains such as cyber-physical systems [3]–[5], biological systems [6], and chemical reaction networks [7]. Despite their wide applicability, however, a theoretical framework that provides formal guarantees as well as scalable computational algorithms for analysis and control of switched stochastic systems remains a major challenge. In particular, the need for such a framework is of immediate importance for safety-critical applications such as autonomous cars and air traffic control [8]. In this work, we target this challenge for switched diffusions, a class of stochastic models where the dynamics are linear and continuous in both time and space domains.

Formal approaches to verification and synthesis for stochastic processes have been the focus of many studies in recent years [9]–[18]. The popularity of these methods is rooted in the formal guarantees that they provide over expressive and succinct specification languages, such as linear temporal logic (LTL), probabilistic computation tree logic (PCTL), and continuous stochastic logic (CSL) [19]. The classical accompanying challenge, though, is the state explosion problem, which is particularly exacerbated for systems with continuous domains. To overcome this problem, approaches based on finite abstractions, which are essentially coarse representations of the process, are proposed [9]–[13]. Existing works, by and large, focus on discrete-time, continuous-space stochastic processes and construct an abstraction in the form of a discrete-time, discrete-space Markov process. This usually results in discrepancies between the abstraction and original system, which are then captured through error bounds. In practice, there are two major limitations to these approaches. One is that most real-world systems evolve in continuous time, and the other is the (lack of) scalability of their computational frameworks. That is, on the one hand, the error bounds can be conservative if the abstraction is not fine enough, and on the other hand fine abstractions can result in state-space explosion.

In this manuscript, we introduce a theoretical and computational synthesis framework for safety properties of (continuous time and space) switched diffusions that is both formal and scalable. The framework consists of two stages: abstraction and synthesis. In the first stage, an appropriate discrete abstraction in the form of an uncertain Markov model that captures all possible behaviors of the system is constructed. This is achieved through a discretization of time and space domains, each introducing an error. These errors are formally characterized and represented as uncertain transition probabilities in the abstraction model. For the space domain, the framework in particular uses a suitable discretization based on the dynamics of the system that results in closed-form analytical solutions for the error term, leading to fast computations. In the second stage, a robust strategy that optimizes a safety property is synthesized over the abstraction. This strategy is computed by considering only the feasible transition probability distributions, preventing the explosion of the error term and resulting in achievable bounds for the safety probability. Finally, this robust strategy is correctly mapped to a switching strategy for...
the system of switched diffusions with the guarantee that the computed safety probability bounds also hold for this system.

The contributions of this work are fourfold. Firstly, we introduce a theoretical framework for formal analysis of continuous-time, continuous-space diffusion processes. This includes the characterization and derivation of the error bounds that result from time discretization. Secondly, we propose a novel space discretization technique that is dynamics-dependent and derive the closed-form analytical solution for the computation of the exact (achievable) error bounds. Thirdly, we introduce a strategy synthesis algorithm for uncertain Markov processes with safety (invariance) properties, which is robust against the embedded uncertainty. Fourthly, we perform an empirical analysis on the trade-off between time and space discretization parameters. One of the main outcomes of this work is a computational framework that is both formal and scalable for switched stochastic processes. The choice of the abstraction model and the derivation of discretization methods with tight error bounds (exact for space) leads to fine and compact abstractions, whose computation is fast without the need to rely on sampling-based approaches (as is currently practiced in the literature, e.g., [12]). Furthermore, due to duality between probabilistic safety and reachability problems [20], the proposed framework can be easily extended to verification and synthesis for more complex properties [14], [21], [22] expressed in, e.g., PCTL and CSL.

II. PROBLEM FORMULATION AND APPROACH

We consider a switched stochastic process that is continuous both in time and space and evolves according to:

$$\begin{align*}
\frac{dx(t)}{dt} &= F(a) x(t) dt + G(a) dw(t), \\
x \in \mathbb{R}^m, \quad a \in A, \quad t \in \mathbb{R}_{\geq 0},
\end{align*}$$

(1)

where

- \( A = \{ a_1, ..., a_{|A|} \} \) is a finite set of actions,
- \( F : A \rightarrow \mathbb{R}^{m \times m} \) is the drift term,
- \( G : A \rightarrow \mathbb{R}^{m \times r} \) is the diffusion term, and
- \( w \) is an \( r \)-dimensional Wiener process.

Under each action \( a \in A \), process \( x \) evolves according to the SDE in (1). An action change at time \( t \) causes a switch in the drift and diffusion values, resulting in an update in the dynamics of the process only, i.e., no change in the value of \( x(t) \) (unlike general stochastic hybrid systems [23]). Then, the process continues evolving according to the updated SDE from \( x(t) \) until the next switch.

**Assumption 1** (Weakly controllable). Under each action \( a \in A \), process (1) is weakly controllable, i.e., the controllability subspace associated to \( (F(a), G(a)) \) is full rank.

This ensures that, under each action \( a \), \( x(t) \) is a non-degenerate Gaussian random variable with values in \( \mathbb{R}^m \) [24].

A sample path or trajectory \( \omega_X \) of \( x \) is a time-unbounded execution of the process (1), i.e., a function \( \omega_X : \mathbb{R}_{\geq 0} \rightarrow \mathbb{R}^m \). We denote the set of all sample paths by \( \Omega \), the segment of a path \( \omega_X \) that is limited to the time interval \([0, t] \) by \( \omega_X^t \), and the value of path \( \omega_X \) at time \( t \) by \( \omega_X(t) \). Let \( F^w_\tau \) be the sigma-algebra generated by the Wiener process \( w(t') \) for \( t' \leq t \) [25]. A switching strategy \( \sigma_X \) is a \( F^w_\tau \)-measurable function that selects actions for the process in (1). Under \( \sigma_X \), the stochastic process \( x \) is defined on the filtered probability space \((\Omega, \mathcal{F}, \mathbb{P}, \text{Prob})\), where \( \mathcal{F} = (\mathcal{F}_t)_{t \geq 0} \) is a filtration such that \( \mathcal{F}_t \subseteq F^w_\tau \), and \( \text{Prob} \) is a probability measure [26], [27].

In this work, we assume that \( \sigma_X \) is also \( \mathcal{F}_t \)-adapted, i.e., its choice of action at time \( t \) depends on \( \omega_X^t \) (the trajectory of \( x \) up to time \( t \)) [26]. In order to avoid Zeno behavior (an infinite number of action switches over a finite time horizon), we also assume that \( \sigma_X \) assigns a finite number of switches over any finite time interval almost surely. Under such a strategy, both the existence and the uniqueness of a solution \( x \) are guaranteed since (1) is a linear SDE. Moreover, \( x \) remains bounded almost surely over a finite time interval [25], [27].

A. Problem Formulation

We are interested in the safety analysis of process \( x \), and our aim is to compute a switching strategy that maximizes the probability of remaining in a safe set.

**Problem 1.** Given the switched stochastic process in (1), a compact and measurable safe set \( X_{\text{safe}} \subseteq \mathbb{R}^m \), and a time duration \( \tau \in \mathbb{R}_{\geq 0} \), find a switching strategy \( \sigma_X \) that maximizes the safety probability of process \( x \), namely the probability of remaining in \( X_{\text{safe}} \) given by

$$P_{\text{safe}}(x, X_{\text{safe}}, \tau \mid \sigma_X) = \text{Prob}(\omega_X(t) \in X_{\text{safe}} \, \forall t \in [0, \tau] \mid \sigma_X).$$

(2)

Safety analysis is one of the fundamental problems in quantitative verification, which can be reformulated as a non-trivial stochastic optimal control problem with a multiplicative cost comprising indicator functions [20], and in particular it has been studied for stochastic hybrid processes [28]. In this paper, the focus is on such analysis, as formulated in Problem 1. It is worth noting that solving this problem can also lead to the solution of the dual problems of reachability analysis [20], again a core problem in formal verification [19]. Therefore, the solution to the safety Problem 1 can be extended to the verification of stochastic hybrid systems over more general PCTL and CSL properties [14], [21], [22].

The measurability of the event in (2) is based on the fact that \( x \) is almost surely continuous. This implies that \( x \) is a separable stochastic process [29, Theorem 38.1], and for separable stochastic processes, (2) is well defined. We emphasize that the measurability of (2) has been established also for more general classes of stochastic processes [30] and that evidently (2) does not account for sets of measures of path zero that might escape the safe set.

B. Proposed Approach

Analytical approaches to Problem 1 are generally infeasible. That is due to the fact that such approaches require solutions to partial differential equations (PDE) with absorbing boundaries, which usually cannot be obtained in closed form [31]. The switched process in (1) can be formulated as a bilinear stochastic control problem with control taking values over a finite set and acting on both drift and diffusion terms.
[32]; however, literature mostly considers actions only in the drift term and studies exclusively classical stochastic control problems, which furthermore require, in general, numerical solutions with an error that is very difficult (if possible at all) to compute [26], [33].

In this work, we take a general approach that is both formal and computationally tractable through a discrete abstraction. We construct a finite model in the form of an uncertain Markov process that captures all possible behaviors of the process in (1). This construction involves a discretization of both time and space domains, which results in a model that approximates the behavior of the continuous process in (1). We quantify the error of this approximation and represent it as uncertainty in the Markov model. We then synthesize an optimal strategy over time. Let \( \tau_{\text{safe}} \) be the sets of paths of \( \mathcal{H} \) that assigns a discrete mode to each path up to time \( t \).

**Assumption 2 (Piecewise-constant \( \sigma_{\mathcal{H}} \)).** The switching strategy \( \sigma_{\mathcal{H}} \) is a piecewise-constant function that may change its value only at time instants \( t = k \cdot \Delta t \), where \( k \in \mathbb{N} \) and \( \Delta t \in \mathbb{R}_{>0} \).

This assumption considers \( \sigma_{\mathcal{H}} \) to be a piecewise-constant function with constant sampling times. Hence, \( \sigma_{\mathcal{H}} \) is a special case of the strategy \( \sigma_X \) defined in Sec. II.

**Definition 2 (SHS execution).** An execution of a SHS \( \mathcal{H} = (A, m, F, G) \) under a switching strategy \( \sigma_{\mathcal{H}} \) and an initial state \( s(0) = (a(0), x(0)) \in A \times \mathbb{R}^m \) is a stochastic process \( s(t) \), whose sample paths for the duration of \( t \in [0, t_f] \) denoted by \( \omega_{\mathcal{H}}^{t_f} \in \text{Paths}_{\mathcal{H}}^{t_f} \) are obtained according to the following algorithm:

1. set \( t = 0 \) and \( \omega_{\mathcal{H}}^{0} = (a(0), x(0)) \);
2. while \( t < t_f \)
   a. \( \sigma = \sigma_{\mathcal{H}}(\omega_{\mathcal{H}}^{t}) \);
   b. compute \( \omega_{\mathcal{H}}^{t} \) according to (1) with initial state \( x(t) \);
   c. extend \( \omega_{\mathcal{H}}^{t} \) by \( (a, \omega_{\mathcal{H}}^{t}) \);
   d. \( t = t + \Delta t \);
end

Given \( \sigma_{\mathcal{H}} \), the probability space defined on process \( x \) naturally extends to process \( s \). We associate to \( x(t) \) a controlled transition probability measure \( T^d \). For \( a \in A \), let \( \mathcal{B}(\mathbb{R}^m) \) be the Borel sigma-algebra on the state space \( \mathbb{R}^m \). Then, starting from a continuous state \( x \in \mathbb{R}^m \) in mode \( a \) at time \( t \in \mathbb{R}_{\geq 0} \), the probability that \( x(t + \Delta t) \in B \), where \( B \in \mathcal{B}(\mathbb{R}^m) \) is a region in the state space \( \mathbb{R}^m \), with the assumption of no mode change in the duration \( \Delta t \) is given by

\[
T^d(B \mid x, a, \Delta t) = \int_B \mathcal{N}(z \mid E_x(\Delta t), Cov_x(\Delta t)) dz, \tag{3}
\]

where \( \mathcal{N}(\cdot \mid E_x(\Delta t), Cov_x(\Delta t)) \) is a multivariate normal density function with mean \( E_x(\Delta t) \) and covariance matrix \( Cov_x(\Delta t) \) given by

\[
E_x(\Delta t) = e^{F(a)\Delta t} x, \tag{4}
\]

\[
Cov_x(\Delta t) = \int_0^{\Delta t} e^{F(a)(\Delta t - z)} G(a)G^T(a) e^{F(a)(\Delta t - z)} dz. \tag{5}
\]

**Lemma 1.** The covariance matrix \( Cov_x(\Delta t) \) is full rank.

Proof. The proof follows directly from Assumption 1. \( \square \)

**B. Finite-state Markov Models**

We utilize Markov models as abstraction structures.

**Definition 3 (MDP).** A Markov Decision Process (MDP) is a tuple \( \mathcal{M} = (Q, A, P) \), where:

- \( Q \) is a finite set of states,
- \( A \) is a finite set of actions,
- \( P : Q \times A \times Q \rightarrow [0, 1] \) is a transition probability function.
The set of actions available at \( q \in Q \) is denoted by \( A(q) \). The function \( P \) has the property \( \sum_{a' \in A} P(q,a,q') = 1 \) for all \( q \in Q \) paired with each \( a \in A(q) \). A path \( \omega \) through an MDP is a sequence of states \( \omega = q_0 \xrightarrow{a_0} q_1 \xrightarrow{a_1} q_2 \xrightarrow{a_2} \ldots \), such that \( a_i \in A(q_i) \) and \( P(q_i,a_i,q_{i+1}) > 0 \) for all \( i \in \mathbb{N} \). We denote the last state of a finite path \( \omega_{\text{fin}} \) by \( \text{last}(\omega_{\text{fin}}) \) and the set of all finite and infinite paths by \( \text{Paths}_{\text{fin}} \) and \( \text{Paths} \), respectively.

A strategy defines a choice of action at each state of the MDP. Its formal definition follows.

**Definition 4** (Strategy). A strategy \( \sigma \) of an MDP model \( M \) is a function mapping a finite path \( \omega_{\text{fin}} = q_0q_1 \ldots q_n \) of \( M \) onto an action in \( A \). In other words, a strategy is a function \( \sigma : \text{Paths}_{\text{fin}} \to A \) that specifies, for every finite path, the next action to be applied. If a strategy depends only on \( \text{last}(\omega_{\text{fin}}) \) and step time \( n \), it is called a (time-dependent) Markov strategy. If a strategy depends only on \( \text{last}(\omega_{\text{fin}}) \), it is called a memoryless or stationary Markov strategy. The set of all strategies is denoted by \( \Sigma \).

Given strategy \( \sigma \), a probability measure \( \text{Prob} \) over the set of all paths \( \text{Paths} \) is induced on the resulting Markov chain [19].

When modeling with MDPs, it might be difficult to determine exact values of transition probabilities between states. In such cases, an interval for each value may be considered. The model that allows the inclusion of these intervals is known as bounded-parameter [35] or interval MDP (IMDP) [22].

**Definition 5** (IMDP). An interval Markov decision process (IMDP) is a tuple \( I = (Q, A, \hat{P}, \hat{P}) \), where \( Q \) and \( A \) are as in Def. 3, and

\[
\begin{align*}
\hat{P} & : Q \times A \times Q \to [0, 1] \text{ is a function, where } \hat{P}(q,a,q') \text{ is the infimum (lower bound) of the transition probabilities from state } q \text{ to state } q' \text{ under action } a \in A(q), \\
\hat{P} & : Q \times A \times Q \to [0, 1] \text{ is a function, where } \hat{P}(q,a,q') \text{ is the supremum (upper bound) of the transition probabilities from state } q \text{ to state } q' \text{ under action } a \in A(q).
\end{align*}
\]

For all \( q, q' \in Q \) and \( a \in A(q) \), it holds that \( \hat{P}(q,a,q') \leq \hat{P}(q,a,q') \) and \( \sum_{q' \in Q} \hat{P}(q,a,q') \leq 1 \leq \sum_{q' \in Q} \hat{P}(q,a,q') \).

Let \( D(Q) \) denote the set of discrete probability distributions over \( Q \). Given \( q \in Q \) and \( a \in A(q) \), we call \( \gamma^a_q = D(Q) \) a feasible distribution reachable from \( q \) by \( a \) if \( \hat{P}(q,a,q') \leq \gamma^a_q(q') \leq \hat{P}(q,a,q') \) for each state \( q' \in Q \). We denote the set of all feasible distributions for state \( q \) and action \( a \) by \( \Gamma^a_q \). In IMDPs, the notions of paths and strategies are extended from MDPs in a straightforward manner. The additional concept is that of adversary, which makes the choice of a feasible distribution.

**Definition 6** (Adversary). Given an IMDP \( I \), an adversary is a function \( \pi : \text{Paths}_{\text{fin}} \times A \to D(Q) \) that, for each finite path \( \omega_{\text{fin}} \in \text{Paths}_{\text{fin}} \) and action \( a \in \text{last}(\omega_{\text{fin}}) \), assigns a feasible distribution \( \pi(\omega_{\text{fin}}, a) \in \Gamma^a_{\text{last}(\omega_{\text{fin}})} \). The set of all adversaries is denoted by \( \Pi \).

Given a finite path \( \omega_{\text{fin}} \), a strategy \( \sigma \), and an adversary \( \pi \), the IMDP evolution proceeds as follows. At state \( q = \text{last}(\omega_{\text{fin}}) \), first an action \( a \in A(q) \) is chosen by \( \sigma \). Then, \( \pi \) resolves the uncertainties and chooses one feasible distribution \( \gamma^a_q \in \Gamma^a_q \). Finally, the next state \( q' \) is chosen according to the distribution \( \gamma^a_q \), and the path \( \omega_{\text{fin}} \) is extended by \( q' \).

Given a strategy \( \sigma \) and an adversary \( \pi \), a probability measure \( \text{Prob}_{\text{fin}} \) over the set of all paths \( \text{Paths} \) (under \( \sigma \) and \( \pi \)) is induced by the resulting Markov chain [12].

### C. Polytopes and their Post Images

Let \( m \in \mathbb{N} \) and consider the \( m \)-dimensional Euclidean space \( \mathbb{R}^m \). A full dimensional polytope \( P \) is defined as the convex hull of at least \( m + 1 \) affinely independent points in \( \mathbb{R}^m \) [36]. The set of vertices of \( P \) is the set of points \( v_1, \ldots, v_{n_P} \in \mathbb{R}^m, n_P \geq m + 1 \), whose convex hull gives \( P \) and with the property that, for any \( i = 1, \ldots, n_P \), point \( v_i \) is not in the convex hull of the remaining points \( v_{i+1}, \ldots, v_{n_P} \). A polytope is completely described by its set of vertices,

\[
P = \text{conv}(v_1, \ldots, v_{n_P}).
\]

where \( \text{conv} \) denotes the convex hull. Alternatively, \( P \) can be described as the bounded intersection of at least \( m + 1 \) closed half spaces. In other words, there exists a \( k \geq m + 1, h_i \in \mathbb{R}^m \), and \( l_i \in \mathbb{R}, i = 1, \ldots, k \) such that

\[
P = \{ x \in \mathbb{R}^m \mid h_i^T x \leq l_i, i = 1, \ldots, k \}.
\]

The above definition can be written as the matrix inequality \( Hx \leq L \), where \( H \in \mathbb{R}^{k \times m} \) and \( L \in \mathbb{R}^k \). Forms (6) and (7) are referred to as \( V \)- and \( H \)-representations of the polytope.

Given a matrix \( M \in \mathbb{R}^{m \times m} \), the post image of polytope \( P \) by \( M \) is defined as [12]:

\[
\text{Post}(P|M) = \{ Mx \mid x \in P \}.
\]

This post image is a polytope itself under the linear transformation \( M \) and can be computed as \( \text{Post}(P|M) = \text{conv}(\{Mv_i, 1 \leq i \leq n_P\}) \).

### IV. DISCRETE ABSTRACTIONS

In the first step of our approach to Problem 1, we model the process in (1) as a SHS and then construct a discrete abstraction of it in the form of an IMDP. This abstraction is based on discretization of both time and space domains. That is, we sample time at a fixed time interval \( \Delta t > 0 \) and partition the given safe set into polytopic regions. Such a discretization, however, introduces an error, which needs to be accounted for as detailed below.

#### A. Stochastic Hybrid System Modeling

We model the switched stochastic process in (1) as a SHS \( H \) as defined in Section III-A. Each action of the process in (1) is associated with a mode of \( H \), and the drift and diffusion terms in mode \( a \) are \( F(a) \) and \( G(a) \), respectively. In this model, each
mode of \( \mathcal{H} \) contains a copy of the safe set \( X_{\text{safe}} \subset \mathbb{R}^m \), and hence, the hybrid safe set is

\[
S_{\text{safe}} = \{ (a, x) \mid x \in X_{\text{safe}} \text{ and } a \in A \}.
\]

Therefore, the safety probability of process \( x(t) \) in (2) becomes

\[
P_{\text{safe}}(s, S_{\text{safe}}, \tau \mid \sigma_\mathcal{H}) = \text{Prob}(\omega_\mathcal{H}(t) \in S_{\text{safe}} \forall t \in [0, \tau] \mid \sigma_\mathcal{H}).
\]

for the process \( s(t) \) of the hybrid system \( \mathcal{H} \), and Problem 1 becomes equivalent to finding \( \sigma_\mathcal{H} \) that maximizes (9).

Figure 1 illustrates a hybrid system representation of a switched stochastic process with two actions. The hybrid system has two modes denoted by \( a_1 \) and \( a_2 \). The black surfaces are the copies of set \( X_{\text{safe}} \) in each mode, and the blue continuous trajectory is a sample path of the hybrid system with two switches. Note that a segment of the path in mode \( a_1 \) exits and re-enters the safe set. Lastly, the values of \( x(t) \) on the trajectory do not change at the times of switching (illustrated as dashed vertical lines); only the modes \( a(t) \) change.

### B. IMDP Abstraction

We abstract the hybrid system \( \mathcal{H} \) to an IMDP \( \mathcal{I} = (Q, A, P, P) \) with the purpose of safety analysis. To this end, we perform a discretization of the hybrid state space \( S \) by distinguishing between the safe and unsafe states and exclusively focusing on \( S_{\text{safe}} \) as described below.

For each mode \( a \in A \), we partition the corresponding safe set \( X_{\text{safe}} \) into a set of cells (regions) that are non-overlapping except for at most trivial sets of measure zero (their boundaries). We assume that each region is a bounded polytope. We denote by \( Q^a = \{ q_1^a, \ldots, q_{|Q^a|}^a \} \) the resulting set of regions in mode \( a \). To each cell \( q_i^a \), we associate a state of the IMDP \( \mathcal{I} \). We overload the notation by using \( q_i^a \) for both a region in \( S \), i.e., \( q_i^a \subseteq \{ a \} \times X_{\text{safe}} \subseteq S_{\text{safe}} \subseteq S \), and a state of \( \mathcal{I} \), i.e., \( q_i^a \in Q \); the exact meaning of it should be clear from the context. Therefore, \( S_{\text{safe}} \) can be represented by \( Q_{\text{safe}} = \bigcup_{a \in A} Q^a \). We define the set of IMDP states to be

\[
Q = Q_{\text{safe}} \cup \{ q_{\text{unsafe}} \},
\]

where \( q_{\text{unsafe}} \) is an IMDP state that represents the remaining hybrid states in \( S \setminus Q_{\text{safe}} \).

We define the set of actions of \( \mathcal{I} \) to be the set of modes \( A \) of \( \mathcal{H} \) and allow all actions to be available at each state of \( \mathcal{I} \), i.e., \( A(q) = A \) for all \( q \in Q \). To capture the safe behavior of \( \mathcal{H} \) by \( \mathcal{I} \), we define the one-step safe transition probability from a continuous state \( x \in X_{\text{safe}} \) to region \( q \in Q_{\text{safe}} \) under action (mode) \( a \in A \) to be

\[
P_{\text{safe}}(q \mid x, a, \Delta t) = \text{Prob}(x(\Delta t) \in q \land \forall t \in [0, \Delta t] (a(t), x(t)) \in S_{\text{safe}} \mid x(0) = x, \quad a(t) = a \forall t \in [0, \Delta t]).
\]

In other words, \( P_{\text{safe}}(q \mid x, a, \Delta t) \) is the probability of reaching \( q \) from \( x \) in \( \Delta t \) without leaving the safe set. By the application of the law of conditional probabilities we get:

\[
P_{\text{safe}}(q \mid x, a, \Delta t) = T^d(q \mid x, a, \Delta t) T^c(\Delta t \mid x, q, a),
\]

where

\[
T^d(q \mid x, a, \Delta t) = \text{Prob}(x(\Delta t) \in q \mid x(0) = x, \quad a(t) = a \forall t \in [0, \Delta t])
\]

is the discrete transition kernel that determines the probability of ending up in \( q \) after \( \Delta t \), as introduced in (3), and

\[
T^c(\Delta t \mid x, q, a) = \text{Prob}((a(t), x(t)) \in S_{\text{safe}} \forall t \in [0, \Delta t] \mid x(0) = x, x(\Delta t) \in q, a(t) = a \forall t \in [0, \Delta t])
\]

is the continuous transition kernel that considers the safety of the continuous trajectories over the time interval \( \Delta t \). Note that \( T^d \) considers all the paths that end up in \( q \) from \( x \). This can be viewed as an error caused by sampling time at intervals of \( \Delta t \), i.e., a time discretization error. The term \( T^c \) corrects this error by considering those paths that exit and then re-enter the safe set during the sampling time interval. As shown in Section IV-D, the computation of \( T^c \) reduces to quantify excursion probabilities for Gaussian processes, which, in general, cannot be obtained analytically [37]. Hence, in what follows we will derive sound upper and lower bounds for \( T^c \).

We define the transitions in \( \mathcal{I} \) by using the safe transition probabilities in (11). The caveat is that the states of \( \mathcal{I} \) are regions in \( \mathcal{H} \), and there are uncountably many possible (continuous) initial states in each region, giving rise to a range of safe feasible transition probabilities to the other regions. Therefore, the exact transition probability from one region to another cannot be known, but its range is given by min and max of (11) over all the possible points in the initial region. That is, for \( q_i, q_j \in Q_{\text{safe}} \) and \( a \in A \), the safe feasible transition probability is bounded from below by

\[
\gamma_{q_i}(q_j) \geq \min_{x \in q_i} P_{\text{safe}}(q_j \mid x, a, \Delta t) = \min_{x \in q_i} T^d(q_j \mid x, a, \Delta t) T^c(\Delta t \mid x, q_j, a)
\]

(12)

\[
\geq \left( \min_{x \in q_i} T^d(q_j \mid x, a, \Delta t) \right) \left( \min_{x \in q_i} T^c(\Delta t \mid x, q_j, a) \right),
\]

(13)
and from above by

\[
\gamma^a_{q_i}(q_j) \leq \max_{x \in q_i} P_{\text{safe}}(q_j \mid x, a, \Delta t) = \max_{x \in q_i} T^d(q_j \mid x, a, \Delta t) T^c(\Delta t \mid x, q_j, a) \leq \left( \max_{x \in q_i} T^d(q_j \mid x, a, \Delta t) \right) \left( \max_{x \in q_i} T^c(\Delta t \mid x, q_j, a) \right),
\]

\( \text{(14)} \quad \text{(15)} \)

Therefore, we can define the transition probability bounds \( \hat{P} \) and \( \hat{P} \) of \( \mathcal{I} \) according to these bounds. Ideally, they should be the exact bounds given in (12) and (14), but they are usually difficult to compute as discussed in Sec. IV-D. Instead, we set the values of \( \hat{P} \) and \( \hat{P} \) according to the possibly looser bounds in (13) and (15), i.e.,

\[
\hat{P}(q_i, a, q_j) = \left( \min_{x \in q_i} T^d(q_j \mid x, a, \Delta t) \right) \left( \min_{x \in q_i} T^c(\Delta t \mid x, q_j, a) \right),
\]

\( \text{(16)} \)

\[
\hat{P}(q_i, a, q_j) = \left( \max_{x \in q_i} T^d(q_j \mid x, a, \Delta t) \right) \left( \max_{x \in q_i} T^c(\Delta t \mid x, q_j, a) \right),
\]

\( \text{(17)} \)

for all \( a \in A \) and \( q_i, q_j \in Q_{\text{safe}} \).

Similarly, we define the bounds of the feasible transition probabilities to the unsafe state as

\[
\gamma^a_{q_i}(q_{\text{unsafe}}) \geq 1 - \max_{x \in q_i} P_{\text{safe}}(X_{\text{safe}} \mid x, a, \Delta t),
\]

\( \text{(18)} \)

\[
\gamma^a_{q_i}(q_{\text{unsafe}}) \leq 1 - \min_{x \in q_i} P_{\text{safe}}(X_{\text{safe}} \mid x, a, \Delta t),
\]

\( \text{(19)} \)

and set the bounds in \( \mathcal{I} \) to be

\[
\hat{P}(q_i, a, q_{\text{unsafe}}) = 1 - \max_{x \in q_i} P_{\text{safe}}(X_{\text{safe}} \mid x, a, \Delta t),
\]

\( \text{(20)} \)

\[
\hat{P}(q_i, a, q_{\text{unsafe}}) = 1 - \min_{x \in q_i} P_{\text{safe}}(X_{\text{safe}} \mid x, a, \Delta t),
\]

\( \text{(21)} \)

for all \( a \in A \) and \( q_i \in Q_{\text{safe}} \). Finally, we make the unsafe state \( q_{\text{unsafe}} \) absorbing, i.e.,

\[
\hat{P}(q_{\text{unsafe}}, a, q_{\text{unsafe}}) = \hat{P}(q_{\text{unsafe}}, a, q_{\text{unsafe}}) = 1,
\]

\( \text{for all } a \in A. \)

Note that the ranges of transition probabilities can be viewed as errors caused by the space discretization. The benefit of using IMDP as the abstraction model is that its structure allows one to encode both time and space discretization errors as uncertainty into the abstraction. In the next three subsections, we show an efficient method of computation for \( \hat{P} \) and \( \hat{P} \), and in Section V, we show how a safe strategy that is robust against the uncertainty can be synthesized over \( \mathcal{I} \). Finally, we prove that the synthesized strategy on \( \mathcal{I} \) can be soundly mapped onto the process in (1) in Section VI.

C. Bounds on the Discrete Transition Kernel

Here, we focus on the values of

\[
\min_{x \in q_i} T^d(q_j \mid x, a, \Delta t), \quad \max_{x \in q_i} T^d(q_j \mid x, a, \Delta t),
\]

\( \text{(22)} \)

and introduce an efficient method for their exact computations. To this end, we first define a hyper-rectangle in \( \mathbb{R}^m \) to be an \( m \)-dimensional rectangle defined by the intervals

\[
[v_{l_1}^{(1)}, v_{u_1}^{(1)}] \times [v_{l_2}^{(2)}, v_{u_2}^{(2)}] \times \cdots \times [v_{l_m}^{(m)}, v_{u_m}^{(m)}],
\]

where vectors \( v_l, v_u \in \mathbb{R}^m \) capture the lower and upper values of the vertices of the rectangle in each dimension, and \( v^{(i)} \) denotes the \( i \)-th component of vector \( v \). Then, we characterize \( T^d \) analytically as follows.

**Proposition 1.** For process \( x \) under action \( a \in A \) for a duration \( \Delta t \), let \( \mathcal{T}_a = \Lambda_a^{-\frac{1}{2}} V_a^T \) be a transformation function (matrix), where \( \Lambda_a = V_a^T \text{Cov}_x(\Delta t) V_a \) is a diagonal matrix whose entries are eigenvalues of \( \text{Cov}_x(\Delta t) \) and \( V_a \) is the corresponding orthonormal (eigenvector) matrix. For a polytopic region \( q \subset \mathbb{R}^m \), if \( \text{Post}(q \mid \mathcal{T}_a) \) is a hyper-rectangle given by \([v_l^{(1)}, v_u^{(1)}] \times \cdots \times [v_l^{(m)}, v_u^{(m)}] \), then it holds that

\[
T^d(q \mid x, a, \Delta t) = \frac{1}{2^m} \prod_{i=1}^m \left( \text{erf}\left( \frac{y^{(i)} - v_l^{(i)}}{\sqrt{2}} \right) - \text{erf}\left( \frac{y^{(i)} - v_u^{(i)}}{\sqrt{2}} \right) \right),
\]

\( \text{(23)} \)

where \( \text{erf}(\cdot) \) is the error function, and \( y^{(i)} \) is the \( i \)-th component of vector \( y = \mathcal{T}_a E_x(\Delta t) \).

**Proof.** For a fixed \( a \in A \) and duration \( \Delta t \), recall that \( T^d(q \mid x, a, \Delta t) \) is given by (3). By applying a whitening transformation through the transformation matrix \( \mathcal{T}_a = \Lambda_a^{-\frac{1}{2}} V_a^T \), we obtain that \( \mathcal{T}_a \text{Cov}_x(\Delta t) \mathcal{T}_a^T = \mathbf{I} \), where \( \mathbf{I} \) is the identity matrix. Thus, by working in the transformed space induced by \( \mathcal{T}_a \), we obtain

\[
T^d(q \mid x, a, \Delta t) = \int_{\text{Post}(q \mid \mathcal{T}_a)} \mathcal{N}(z \mid \mathcal{T}_a E_x(\Delta t), \mathbf{I}) \, dz.
\]

Under the assumption that \( \text{Post}(q \mid \mathcal{T}_a) \) is a hyper-rectangle, the above multidimensional integral can be separated and expressed as a product of \( m \) integrals of uni-dimensional normal distributions:

\[
T^d(q \mid x, a, \Delta t) = \int_{\text{Post}(q \mid \mathcal{T}_a)} \mathcal{N}(z \mid \mathcal{T}_a E_x(\Delta t), \mathbf{I}) \, dz
= \int_{v_l^{(1)}}^{v_u^{(1)}} \cdots \int_{v_l^{(m)}}^{v_u^{(m)}} \mathcal{N}(z_1 \mid y^{(1)}, 1) \cdots \mathcal{N}(z_m \mid y^{(m)}, 1) \, dz_1 \cdots dz_m
= \prod_{i=1}^m \frac{1}{2} \left( \text{erf}\left( \frac{y^{(i)} - v_l^{(i)}}{\sqrt{2}} \right) - \text{erf}\left( \frac{y^{(i)} - v_u^{(i)}}{\sqrt{2}} \right) \right),
\]

where \( y = \mathcal{T}_a E_x(\Delta t) \).

\( \Box \)

A direct consequence of Proposition 1 is that the optimization problems in (22) can be performed on (23) through a linear transformation, as stated by the following corollary.

**Corollary 1.** For polytopic regions \( q_i, q_j \subset \mathbb{R}^m \) and process \( x \) under action \( a \) for the duration of \( \Delta t \), assume \( \text{Post}(q_i \mid \mathcal{T}_a) \) is a hyper-rectangle given by

\[
[v_l^{(1)}, v_u^{(1)}] \times \cdots \times [v_l^{(m)}, v_u^{(m)}],
\]

and define \( q'_i = \text{Post}(q_i \mid e^{P(x)\Delta t}) \) and

\[
f(y) = \frac{1}{2^m} \prod_{i=1}^m \left( \text{erf}\left( \frac{y^{(i)} - v_l^{(i)}}{\sqrt{2}} \right) - \text{erf}\left( \frac{y^{(i)} - v_u^{(i)}}{\sqrt{2}} \right) \right),
\]
Then, it holds that
\[
\min_{q \in T_a} T^d(q_j | x, a, \Delta t) = \min_{y \in \text{Post}(q'_i | T_a)} f(y),
\]
\[
\max_{q \in T_a} T^d(q_j | x, a, \Delta t) = \max_{y \in \text{Post}(q'_i | T_a)} f(y).
\]

The above proposition and corollary show that, with a particular geometry for \( q_i \), an analytical form can be obtained for the discrete kernel. This is an important observation because it enables efficient computation for the min and max values of the kernel. Therefore, we use a space discretization to satisfy the condition in Proposition 1, as described below.

For each mode \( a \in A \), we define the linear transformation function (matrix) of
\[
T_a = \Lambda_a^{-\frac{1}{2}} V_a^T,
\]
where \( \Lambda_a = V_a^T \text{Cov}_X(\Delta t) V_a \) is a diagonal matrix whose entries are the eigenvalues of \( \text{Cov}_X(\Delta t) \), and \( V_a \) is the corresponding orthonormal (eigenvector) matrix. The discretization of \( X_{\text{safe}} \) in mode \( a \) is achieved by using a grid in the transformed space by \( T_a \). That is, we first transform \( X_{\text{safe}} \) by \( T_a \), and then discretize it using a grid, and then transform it back to the original space using \( T_a^{-1} \). This method of discretization guarantees that, for each \( q^a \in Q^a \), \( \text{Post}(q^a | T_a) \) is a hyper-rectangle. Hence, we can use the result of Proposition 1 and Corollary 1 for the computation of the values in (22).

It is worth noting that, for an arbitrary geometry of \( X_{\text{safe}} \), it may not be possible to obtain a discretization such that \( \bigcup_{q^a \in Q^a} q^a = X_{\text{safe}} \). Nevertheless, for the purpose of safety analysis, a discretization that under-approximates \( X_{\text{safe}} \), i.e., \( \bigcup_{q^a \in Q^a} q^a \subseteq X_{\text{safe}} \), is sufficient in each mode \( a \). For a better approximation, the grid can be non-uniform, allowing in particular for smaller cells near the boundary of \( X_{\text{safe}} \), as in [11].

In the next theorem, we use the result of Corollary 1 and the Karush-Kuhn-Tucker (KKT) conditions [38] to compute the exact values for (22).

**Theorem 1.** For polytopic regions \( q_i, q_j \subset \mathbb{R}^m \) and transformation matrix \( T_a \), let \( \text{Post}(q_j | T_a) \) be a hyper-rectangle defined by intervals
\[
[v_1^{(i)}_u, v_1^{(i)}_l] \times [v_2^{(i)}_u, v_2^{(i)}_l] \times \cdots \times [v_m^{(i)}_u, v_m^{(i)}_l],
\]
and
\[
\text{Post}(q'_i | T_a) = \{ y \in \mathbb{R}^m | Hy \leq b \},
\]
where \( q'_i = \text{Post}(q_i e^{F(\alpha) \Delta t}) \), \( H \in \mathbb{R}^{k \times m} \), \( b \in \mathbb{R}^m \), and \( k \geq m + 1 \), and call
\[
f(y) = \frac{1}{2m} \prod_{i=1}^m \left( \text{erf}(\frac{y^{(i)}_u - v^{(i)}_l}{\sqrt{2}}) - \text{erf}(\frac{y^{(i)}_l - v^{(i)}_u}{\sqrt{2}}) \right).
\]

Introduce the following conditions:

- **Condition 1:** \( y \) is at the center of \( \text{Post}(q'_i | T_a) \), i.e.,
\[
y = \frac{v^{(1)}_u + v^{(1)}_l}{2}, \ldots, \frac{v^{(m)}_u + v^{(m)}_l}{2}.
\]
- **Condition 2:** \( y \) is a vertex of \( \text{Post}(q'_i | T_a) \).
- **Condition 3:** \( y \) is on the boundary of \( \text{Post}(q'_i | T_a) \), where \( r \geq 1 \) of the \( k \) half-spaces that define \( \text{Post}(q'_i | T_a) \) intersect, and
\[
\nabla f(y) = \bar{H}^T \mu,
\]
for vector \( \mu = (\mu_1, \ldots, \mu_r) \) of non-negative constants, and submatrix \( \bar{H} \in \mathbb{R}^{r \times m} \) that contains only the rows of \( H \) that correspond to the \( r \)-intersecting half-spaces at \( y \).
- **Condition 4:** \( y \) is as in Condition 3, and
\[
\nabla f(y) = -\bar{H}^T \mu,
\]
for vector \( \mu = (\mu_1, \ldots, \mu_r) \) of non-negative constants, and \( H \) is defined as in Condition 3.

Then, it follows that the point \( y \in \text{Post}(q'_i | T_a) \) that satisfies
Condition 1 necessarily maximizes \( f(y) \). If Condition 1 cannot be satisfied, then the maximum is necessarily given by one of the points that satisfy Condition 2 or 3. Furthermore, the point \( y \in \text{Post}(q'_i | T_a) \) that minimizes \( f(y) \) necessarily satisfies Condition 2 or 4.

The proof of Theorem 1 can be found in Appendix A.

Theorem 1 identifies the arguments (points \( y \) in \( \text{Post}(q'_i | T_a) \)) that give rise to the optimal values of the discrete transition kernel in (22). Then, the actual optimal values of \( T^d \) can be computed by (24) as guaranteed by Corollary 1. Therefore, from Theorem 1, an algorithm can be constructed to generate a set of finite candidate points based on Conditions 1-4 and to obtain the exact values of (22) by plugging those points into (24).

In short, Condition 1 maximizes the unconstrained problem and gives rise to the global maximum. Hence, if the center of \( q_i \) is contained in \( \text{Post}(q'_i | T_a) \), no further check is required for maximum. If not, the maximum is given by a point on the boundary of \( \text{Post}(q'_i | T_a) \). If it is either a vertex (Condition 2) or a boundary point that satisfies Condition 3. The minimum is always given by a boundary point, which can be either a vertex or a boundary point that satisfies Condition 4. Note that Conditions 3 and 4 are similar and both state that the optimal value of \( T^d \) is given by a point where the gradient of \( T^d \) becomes linearly dependent on the vectors that are defined by the intersecting half-spaces of \( \text{Post}(q'_i | T_a) \) at that point. Each of these two conditions defines a system of \( m \) equations and \( r < m \) variables, which may have a solution only if some of the equations are linear combinations of the others.

In order to give a better insight on the conditions and the result of Theorem 3, an illustration on a 2-dimensional system is provided in the example below.

**Example 1.** Consider a 2-dimensional system (\( m = 2 \)), where \( \text{Post}(q_j | T_a) \) is a rectangle given by \( \left[ v_1^{(1)}_u, v_1^{(1)}_l \right] \times \left[ v_2^{(2)}_u, v_2^{(2)}_l \right] \), and \( \text{Post}(q'_i | T_a) \) is defined by the following four inequalities for \( y \in \mathbb{R}^2 \):
\[
h_{11} y^{(1)} + h_{12} y^{(2)} \leq b^{(i)}, \quad i \in \{1, 2, 3, 4\}.
\]

Theorem 1 guarantees that if point \( \bar{y} = \left( \frac{y^{(1)} + v^{(1)}_l}{2}, \frac{y^{(2)} + v^{(2)}_l}{2} \right) \)
satisfies all the four inequalities (Condition 1), then
\[
\max_{x \in \mathcal{H}} T^d(q_j | x, a, \Delta t) = f(\bar{y}).
\]
If not, then the maximum is given by a point on the lines
\[
h_{11} y^{(1)} + h_{12} y^{(2)} = b^{(i)}, \quad i \in \{1, 2, 3, 4\}.
\]
That point is either a vertex, i.e., an intersection of two lines, (Condition 2) or where
\[ \frac{\partial f(y)}{\partial y^{(1)}} = h_{i1} \mu_i, \quad \text{and} \quad \frac{\partial f(y)}{\partial y^{(2)}} = h_{i2} \mu_i. \]
(Condition 3), resulting in
\[ \frac{\partial f(y)}{\partial y^{(2)}} = \frac{h_{i2}}{h_{i1}} \frac{\partial f(y)}{\partial y^{(1)}}. \]
Note that \( \frac{h_{i2}}{h_{i1}} \) is the slope of the line perpendicular to the i-th boundary line. This means that a non-vertex candidate point on line i is where the gradient of \( f(y) \) becomes perpendicular to the line. It is a valid candidate if it is on the line segment that satisfies the four inequalities, i.e., on the boundary of Post(\( q_i | T \)). Let \( Y \) be the set containing all such valid candidate points and the 4 vertices, i.e., \( 4 \leq |Y| \leq 8 \). Then,
\[ \max_{x \in Y} T_c(q_j | x, a, \Delta t) = \max_{y \in Y} f(\bar{y}). \]
For the minimum, all four vertices need to be considered (Condition 2) in addition to the valid points where the gradient of \(- f(y)\) becomes perpendicular to the boundary line (Condition 4). Let \( Y' \) denote the set of these points, i.e., \( 4 \leq |Y'| \leq 8 \). Then,
\[ \min_{x \in Y} T_c(q_j | x, a, \Delta t) = \min_{y \in Y} f(\bar{y}). \]
It is important to note that Theorem 1 enables to find the values for \( \overline{\mu} \) without the need for sampling-based methods and/or for over-approximations of the current value, as currently done in the literature, e.g., [1], [9], [12]. Thus, in addition to computing the exact values for those quantities, it also speeds up the process of building the abstraction.

D. Bounds on the Continuous Transition Kernel

Here, we derive bounds for the quantities
\[ \min_{x_1 \in q_i} T_c^c(\Delta t | x_1, q_j, a), \quad \max_{x_1 \in q_i} T_c^c(\Delta t | x_1, q_j, a). \] (25)
To this end, we first need to introduce the notion of bridge of a stochastic process [39]. Intuitively, the bridge of a stochastic process is also a stochastic process for which both initial and final states are known. Formally, for \( \mathbf{x}(t) \) with given initial and final states \( x_1, x_2 \in \mathbb{R}^m \), fixed action \( a \in A \), and duration \( \Delta t \in \mathbb{R}_{\geq 0} \), the bridge \( \mathbf{b}_c^{x_1, x_2, a, \Delta t}(t) \) (or simply \( \mathbf{b}_c(t) \) if the context is clear) is defined as
\[ \mathbf{b}_c^{x_1, x_2, a, \Delta t}(t) = \mathbf{x}(t), \quad \text{conditioned on} \] \[ \mathbf{x}(0) = x_1, \quad \mathbf{x}(\Delta t) = x_2, \quad a(t) = a \quad \forall t \in [0, \Delta t]. \]
The following proposition guarantees that \( \mathbf{b}_c(t) \) is a Gaussian process and derives its expectation and covariance from those of \( \mathbf{x} \) [39], [40].

Proposition 2 (Proposition 4 in [41]). Let \( \mathbf{x} \) be the stochastic process described by (1). For \( a \in A, x_1, x_2 \in \mathbb{R}^m \), \( \mathbf{b}_c^{x_1, x_2, a, \Delta t} \) is a Gaussian process with a probability measure \( \text{Prob}^{b_c} \) such that \( \mathbf{b}_c(t) \sim \mathcal{N}(E_{b_c}(t), \text{Cov}_{b_c}(t)) \), and with the following expectation and covariance
\[ E_{b_c}(t) = \text{Cov}_{b_c}(\Delta t, t) \text{Cov}_{b_c}(\Delta t)^{-1} (x_2 - E_{b_c}(\Delta t)) + E_{b_c}(t), \] (26)
\[ \text{Cov}_{b_c}(t) = \text{Cov}_{b_c}(t) - \text{Cov}_{b_c}(\Delta t, t) \text{Cov}_{b_c}(\Delta t)^{-1} \] (Cov_{b_c}(\Delta t, t), (27)
where \( \text{Cov}_{b_c}(t) \) is the covariance of process \( \mathbf{x} \) at times \( t_1 \) and \( t_2 \) given by
\[ \text{Cov}_{b_c}(t_1, t_2) = \mathbb{E} \left[ (\mathbf{x}(t_1) - E_{b_c}(t_1)) (\mathbf{x}(t_2) - E_{b_c}(t_2))^T \right] = e^{F(a) t_2} \text{Cov}_{b_c}(0) (e^{F(a) t_1})^T + \int_{0}^{\min(t_1, t_2)} e^{F(a) (t_2 - u) G(a) G(a)^T} \left( e^{F(a) (t_1 - u)} \right)^T du. \] (28)
Process \( \mathbf{b}_c(t) \) describes the evolution of \( \mathbf{x} \) during \([0, \Delta t] \) given that \( \mathbf{x}(0) \) and \( \mathbf{x}(\Delta t) \) are known. Since \( \mathbf{b}_c \) is Gaussian, we can use the theory of Gaussian processes to derive bounds for \( T_c^c(\Delta t | x_1, q_j, a) \). Theorem 2 focuses on the upper bound.

Theorem 2. For polytopic regions \( q_i, q_j \subseteq X_{\text{safe}} \), it holds that
\[ \max_{x_1 \in q_i} T_c^c(\Delta t | x_1, q_j, a) \leq \max_{x \in q_i} \int_{X_{\text{safe}}} N(z | \bar{x}, \text{Cov}_{b_c}(\frac{\Delta t}{2})) dz, \] (29)
where
\[ \bar{q} = \text{Post}(q_i | A_1) \oplus \text{Post}(q_j | A_2), \]
\[ A_1 = e^{F(a) \frac{\Delta t}{2}} - \text{Cov}_{b_c}(\Delta t, \frac{\Delta t}{2}) \text{Cov}_{b_c}(\Delta t)^{-1} e^{F(a) \Delta t}, \]
\[ A_2 = \text{Cov}_{b_c}(\Delta t, \frac{\Delta t}{2}) \text{Cov}_{b_c}(\Delta t)^{-1}, \]
and \( \oplus \) is the Minkowski sum.

Proof. Let \( D = \{t_1, t_2, \ldots\} \subseteq [0, \Delta t] \) be a countable and dense subset of \([0, \Delta t] \) such that \( \frac{\Delta t}{2} \notin D \). The upper bound on the continuous kernel from region \( q_i \) to \( q_j \) is the maximum safety probability of the bridge over all the possible initial points in \( q_i \) and final points in \( q_j \):
\[ \max_{x_1 \in q_i} T_c^c(\Delta t | q_j, x_1, a) \]
\[ = \max_{(x_1, x_2) \subseteq (q_i, q_j)} \text{Prob}^{b_c}(\mathbf{b}_c(t_i) \in X_{\text{safe}}, \forall t_i \in D) \]
\[ \leq \max_{(x_1, x_2) \subseteq (q_i, q_j)} \text{Prob}^{b_c}(\mathbf{b}_c^c(\frac{\Delta t}{2}) \in X_{\text{safe}}) \]
\[ = \max_{(x_1, x_2) \subseteq (q_i, q_j)} \int_{X_{\text{safe}}} N(z | E_{b_c}(\frac{\Delta t}{2}), \text{Cov}_{b_c}(\frac{\Delta t}{2})) dz, \]
where in particular the first identity is due to the separability of \( \mathbf{b}_c \) in \([0, \Delta t] \) and the inequality is due to the fact that \( \frac{\Delta t}{2} \notin D \).
In the above optimization, the only term that depends on \( x_1 \) and \( x_2 \) is \( E_{b_c}(\frac{\Delta t}{2}) \). By rearranging the terms in (26), we have
\[ E_{b_c}(\frac{\Delta t}{2}) = \left( e^{F(a) \frac{\Delta t}{2}} - \text{Cov}_{b_c}(\Delta t, \frac{\Delta t}{2}) \text{Cov}_{b_c}(\Delta t)^{-1} e^{F(a) \Delta t} \right) x_1 + \left( \text{Cov}_{b_c}(\Delta t, \frac{\Delta t}{2}) \text{Cov}_{b_c}(\Delta t)^{-1} \right) x_2. \]
Theorem 2 shows that the upper bound of $T^c$ is given by a constrained maximization (over a convex region $q$) of the integral of a normal distribution over $X_{\text{safe}}$. An efficient method of computation for this maximization problem is similar to the one introduced in Section IV-C. That is to solve the problem in the transformed space where $\text{Cov}_{b_{\mathbf{x}}}(\mathbf{A})$ is the identity matrix. Then, through a grid discretization of the Minkowski sum of the linear transformations of $q_i$ and $q_j$, the above optimization can be defined as a maximization over all the points in the $\text{Minkowski} \text{ sum of the linear transformations of } q_i \text{ and } q_j$. □

In order to derive the lower bound of $T^c(\Delta t \mid q_i, x_1, a)$, we employ the following terms need to be defined first.

1. Intuitively, $b^{(i)}_{\mathbf{x}}(t)$ represents the $i$-th component of the bridge with zero mean. Furthermore, for $i \in \{1, \ldots, m\}$, we use $K^{d,i}_{b_{\mathbf{x}}}$ to denote the constant that bounds the expected maximum variation of the process $b^{(i)}_{\mathbf{x}}(t)$.

2. $\epsilon_q = \min_{x \in q} ||x - \partial X_{\text{safe}}||_1$, where $\partial X_{\text{safe}}$ is the boundary of $X_{\text{safe}}$.

3. $E_{b_{\mathbf{x}}}(\mathbf{A})$ is the sum of linear transformations of $x_1$ and $x_2$. Since $x_1 \in q_i$ and $x_2 \in q_j$, the above optimization can be defined as a maximization over all the points in the Minkowski sum of the linear transformations of $q_i$ and $q_j$.

Therefore, $E_{b_{\mathbf{x}}}(\mathbf{A})$ is the identity matrix. Then, through a grid discretization of the Minkowski sum of the linear transformations of $q_i$ and $q_j$, the above optimization can be defined as a maximization over all the points in the Minkowski sum of the linear transformations of $q_i$ and $q_j$.
We can also efficiently compute bounds for these quantities by using the results obtained above. The following proposition shows this efficient method of computation.

**Proposition 3.** Let \( Q^a \) and \( Q^n \) be two sets of polytopic regions in mode \( a \) such that

\[
\bigcup_{q \in Q^a} q \subseteq X_{\text{safe}} \subseteq \bigcup_{q \in Q^n} q,
\]

and \( \text{Post}(q|T_a) \) is a hyper-rectangle defined by

\[
[v_i(1), v_u(1)] \times \cdots \times [v_i(m), v_u(m)]
\]

for every \( q \in Q^a \cup \bar{Q}^n \), and call

\[
f(y, q) = \frac{1}{2m} \prod_{i=1}^{m} \left( \frac{y^{(i)} - v^{(i)}(1)}{\sqrt{2}} - \text{erf} \left( \frac{y^{(i)} - v^{(i)}(1)}{\sqrt{2}} \right) \right).
\]

Then, it holds that

\[
\max_{x \in Q^a} \left[ P_{\text{safe}}(X_{\text{safe}} | x, a, \Delta t) \right] \leq \max_{y \in \text{Post}(q|T_{a})} \sum_{q \in Q^n} f(y, q), \tag{31}
\]

\[
\min_{x \in Q^a} \left[ P_{\text{safe}}(X_{\text{safe}} | x, a, \Delta t) \right] \geq \min_{y \in \text{Post}(q|T_{a})} \sum_{q \in Q^n} f(y, q), \tag{32}
\]

\[
\text{where } q_i = \text{Post}(q_i|e^{f(a)\Delta t}),
\]

\[
\alpha = \begin{cases} 
\max\{0, 1 - 2 \sum_{i=1}^{m} e^{-\frac{\eta_{x,i}}{\pi \epsilon}} \} & \text{if } \eta_{x,i} > 0, \\
0 & \forall i \in \{1, \ldots, m\}
\end{cases}
\]

\[
\eta_{x,i} = \epsilon_{q_i} - \left( L_{x}^{(i)} + 12 \int_{0}^{\frac{1}{2} K_{x}^{d,1} \Delta t} \ln \left( \frac{K_{x}^{d,1} \Delta t}{2z} + 1 \right) dz \right),
\]

\[
\text{and } \epsilon_{q_i}, K_{x}^{d,1} \text{ and } L_{x}^{(i)} \text{ are the constants introduced in Section IV-D, but computed for process } x^{(i)}.
\]

The proof of this proposition is in Appendix C.

Intuitively, Proposition 3 states that, with a particular choice of discretization, i.e., a grid in the transformed space, the safe transition probability to \( X_{\text{safe}} \) is equal to the sum of the transition probabilities to the discrete regions, where each discrete transition kernel is given by the close-form function \( f(y, q) \) in (31). If \( X_{\text{safe}} \) cannot be precisely discretized with a grid (in the transformed space), then the upper and lower bounds of the transition probabilities are given by the over- and under-approximating grids (\( \tilde{Q}^n \) and \( Q^n \)), respectively. In both cases, the bounds need to be scaled by the time discretization error, which is captured by \( \alpha \) in (33) for the lower bound. For the upper bound in (32), the maximum time discretization error is taken to be one. Theoretically, a tighter upper bound can be obtained by extending Theorem 2 for the time discretization error. In the experimental evaluations, however, it made no practical improvement to the bounds whilst drastically increasing the burden of computations.

Note that, for the computation of the values in (32) and (33), a straightforward extension of Theorem 1 can be used. In the extended version, the point \( y \) at which \( \sum_{q \in \bar{Q}^n} \nabla f(y, q) = 0 \) needs to be considered instead of the center point in Condition 1. Condition 2 remains the same, and in Conditions 3 and 4, \( \nabla f(y) \) needs to be replaced with \( \sum_{q \in Q^n} \nabla f(y, q) \) and \( \sum_{q \in \tilde{Q}^n} \nabla f(y, q) \), respectively. This leads to an efficient computation of the transition probability bounds to \( q_{\text{unsafe}} \), finalizing the construction of the abstraction \( I \).

**V. SYNTHESIS**

Given the IMDP abstraction \( I \), our goal is to synthesize a strategy that is robust against all the introduced uncertainties (errors) by the discretization of time and space domains. These uncertainties can be viewed as the nondeterministic choice of a feasible transition probability from one IMDP state to another under a given action. Therefore, we interpret the evolution of the IMDP as a 2-player stochastic game, where Player 1 chooses an action \( a \in A \) at state \( q \in Q \), and Player 2 chooses a feasible transition probability distribution \( \gamma_q^a \in \Gamma_q^a \). This game is adversarial, where the objectives of Players 1 and 2 are to maximize and minimize the probability of remaining in the safe set, respectively. Hence, the goal becomes to synthesize a strategy for Player 1 that is robust against all adversaries.

It is worth noting that, in this 2-player stochastic game, the choice for Player 2 is from a continuous set, and hence the classical algorithms for 2-player games, e.g., [43] cannot be used. In [12], a synthesis algorithm for reachability for IMDPs is introduced. Since reachability and safety problems are dual, one can adapt that algorithm to solve for the safety problem. The dual reachability problem of safety is to find a strategy that minimizes the probability of reaching the unsafe state under all possible adversarial strategies. We introduce a synthesis algorithm that directly solves the safety problem in this section.

Recall that \( \tau \) is the required time duration for the system to remain safe. Let \( k_r = \lceil \frac{T}{\tau} \rceil \) be the equivalent number of time steps, where \( \lceil \cdot \rceil \) is the ceiling function. Also note that, given the IMDP abstraction \( I \), under a strategy \( \sigma \), the probability of remaining safe from each state is necessarily a range for all the available choices of Player 2 or adversaries. Let \( \hat{p}^k_\sigma(q) \) and \( \hat{p}^0_\sigma(q) \) denote the lower and upper bounds of the probability of remaining safe in \( k \) time steps starting from state \( q \in Q \) under strategy \( \sigma \), respectively. Derived from Bellman equation, we can compute the optimal lower bound by \( k_r \) recursive evaluations of

\[
\hat{p}^k_\sigma(q) = \max_{a \in A(q)} \min_{q' \in Q} \sum_{q \in Q} \gamma_q^a(q') \hat{p}^{k-1}_\sigma(q') \tag{34}
\]

with initial values of \( \hat{p}^0_\sigma(q) = 1 \) for \( q \in Q_{\text{safe}} \) and \( \hat{p}^0_\sigma(q_{\text{unsafe}}) = 0 \).

The minimization over the adversaries can be computed iteratively through an ordering of the states of \( I \) [12], [35]. Let \( O^1 = o_1, \ldots, o_{|Q|} \), where \( o_i \in Q \) be an ascending ordering of the states in \( Q \) with respect to the safety probability. Then, the adversary that minimizes the safety probability in one transition is the one that assigns as much transition probability mass as possible to the states early in the ordering \( O^1 \). For the state-action pair \( (q, a) \), let \( r \) be the state index \( 1 \leq r \leq |Q| \).
in the ordering $O^\dagger$ that maximizes the following expression without letting it exceed 1:

$$\sum_{i=1}^{r-1} \hat{P}(q_i, a_i) + \sum_{i=r}^{\vert Q \vert} \hat{P}(q_i, a_i).$$

Then, the minimizing adversary at $(q, a)$ is:

$$\gamma^*_{q}(a) = \begin{cases} \hat{P}(q_i, a_i) & \text{if } i < r \\ \hat{P}(q_r, a_r) & \text{if } i > r \end{cases},$$

$$\gamma^*_{q}(a_r) = 1 - \sum_{i=1, i \neq r}^{\vert Q \vert} \gamma^*_{q}(a_i).$$

Once these adversaries are obtained for all $a \in A(q)$, then a maximization over the actions can be performed to complete the computation for one time step in (34). This operation can be achieved by a matrix-vector multiplication followed by a maximization, as detailed in [12]. After $k_r$ iterations of this algorithm, the robust strategy $\sigma^*$ and the lower bound safety probability $\hat{p}^k_{\sigma^*}(q)$ for each $q \in Q$ are obtained. The upper bounds are then given by recursive evaluations of

$$\hat{p}^k_{\sigma^*}(q) = \max_{\gamma^*_{\sigma^*}(q') \in P^*_{\gamma^*}(q)} \sum_{q' \in Q} \gamma^*_{q}(q') \hat{p}^{k-1}_{\sigma^*}(q'),$$

with the initial values of $\hat{p}^0_{\sigma^*}(q) = 1$ for $q \in Q_{safe}$ and $\hat{p}^0_{\sigma^*}(q_{unsafe}) = 0$. The maximization over adversaries is obtained through (35) with a descending ordering $O^\dagger$ of the states.

The complexity of the above strategy synthesis algorithm is polynomial in the size of the IMDP $\mathcal{I}$, and the obtained strategy is a (time-dependent) Markov strategy for finite $k_r$ [12]. Even though this extension is beyond the scope of this work, it is worth noting that the Bellman equations in (34) and (36) are guaranteed to converge as $k \to \infty$ [12], [44]. Therefore, the safety computations can also be performed for unbounded (safety) time durations, i.e., $\tau \to \infty$, in which case the obtained strategy becomes stationary (memoryless).

VI. Correctness Guarantees and Completeness Analysis

The computed strategy $\sigma^*$ and probability bounds $\hat{p}^k_{\sigma^*}$ and $\hat{p}^k_{\sigma^*}$ on $\mathcal{I}$ also hold for the hybrid system $\mathcal{H}$. Let $\mathcal{L} : Paths_{\mathcal{H}}^\infty \rightarrow Paths_{\mathcal{I}}^\infty$ be a function that maps the sample paths of the hybrid system $\mathcal{H}$ to the finite paths of the IMDP $\mathcal{I}$ through sampling time at intervals $\Delta t$, i.e., $\mathcal{L}(\omega_{k\Delta t}^\mathcal{H}) = q_0, q_1, \ldots, q_k$ for all $k \geq 0$, where $q_i$ is the state in $\mathcal{I}$ that corresponds to the region $q_i^\mathcal{H}(i\Delta t)$ in mode $a(i\Delta t)$ of $\mathcal{H}$ such that $x(i\Delta t) \in q_i^\mathcal{H}(i\Delta t)$ for all $0 \leq i \leq k$. Then, the IMDP strategy $\sigma^*$ correctly maps to a piecewise-constant switching strategy $\sigma^*_{\mathcal{H}}$ of $\mathcal{H}$ through

$$\sigma^*_{\mathcal{H}}(\omega_{k\Delta t}^\mathcal{H}) = \sigma^*_{\mathcal{I}}(\mathcal{L}(\omega_{k\Delta t}^\mathcal{H})).$$

Using this construction of the switching strategy $\sigma^*_{\mathcal{H}}$, the following theorem shows that the safety probability bounds $\hat{p}^k_{\sigma^*}$ and $\hat{p}^k_{\sigma^*}$ are guaranteed to hold for the safety of process $s$ in $\mathcal{H}$ under $\sigma^*_{\mathcal{H}}$ for the duration of $\tau$.

**Theorem 4.** Let $s$ be the execution associated to the hybrid system $\mathcal{H}$ and $\mathcal{I}$ be the IMDP abstraction of $\mathcal{H}$ for sampling time $\Delta t > 0$, as described in Section IV. Furthermore, let $\sigma^*$ be the strategy on $\mathcal{I}$ computed by solving (34) and (36) with safety probability bounds $\hat{p}^k_{\sigma^*}$ and $\hat{p}^k_{\sigma^*}$, where $k_r = \lceil \frac{\tau}{\Delta t} \rceil$ for time duration $\tau \in \mathbb{R}_{\geq 0}$. For switching strategy $\sigma^*_{\mathcal{H}}$ constructed from $\sigma^*$ per (37) and initial hybrid state $s(0) = s_0 = (a_0, x_0)$ with $x_0 \in q_0$ in mode $a_0 \in A$, it holds that

$$P_{safe}(s, S_{safe}, \tau \mid \sigma^*_{\mathcal{I}}, s_0) \in \left[ \hat{p}^k_{\sigma^*}(q_0), \hat{p}^k_{\sigma^*}(q_0) \right],$$

where $P_{safe}(s, S_{safe}, \tau \mid \sigma^*_{\mathcal{I}}, s_0)$ is the safety probability for the process $s$ initialized at $s_0$ under $\sigma^*_{\mathcal{H}}$.

**Proof.** By Theorems 1, 2 and 3, it holds that $\Gamma^a_{q_0}$ contains all the feasible distributions of $s = (a, x)$ during $[0, \Delta t]$ given $a(t) = a$ and $x(0) = x$, for all $x \in q_0$. Bounds $\hat{p}^k_{\sigma^*}(q_0)$ and $\hat{p}^k_{\sigma^*}(q_0)$ are computed by solving (34) and (36). Thus, given that the system is Markov, by induction over all the discrete time intervals we have that $\hat{p}^k_{\sigma^*}(q_0)$ is the lower bound over all possible distributions of $s$ during $[0, \tau]$, and $\hat{p}^k_{\sigma^*}(q_0)$ is the upper bound on the same set.

Theorem 4 guarantees that the safety probability of the process $s$ is contained in the safety probability interval computed on the abstraction $\mathcal{I}$. The size of this interval depends on the difference of the one-step transition probability bounds of $\hat{P}$ and $\hat{P}$ in $\mathcal{I}$, which can be viewed as the error induced by time and space discretization of $\mathcal{H}$ cast into the abstraction $\mathcal{I}$. Both of these errors can be tuned: for time, the discretization variable is $\Delta t$; for space, the discretization variable is the largest volume of the discrete cells (partitions) denoted by $V_{max}$. Similar to the proof in [1] for switching diffusions, it can be shown that, as both $\Delta t \to 0$ and $V_{max} \to 0$, the error of the abstraction goes to 0, and the IMDP abstraction becomes an MDP, i.e., $\hat{P}(q, a, q') \rightarrow P(q, a, q') \leftarrow \hat{P}(q, a, q')$ for all $q, q' \in Q$ and $a \in A(q)$; hence, $\hat{p}^k_{\sigma^*}(q_0) \to P_{safe}(s, S_{safe}, \tau \mid \sigma^*_{\mathcal{H}}, s_0) \leftarrow \hat{p}^k_{\sigma^*}(q_0)$. As a consequence, the synthesized strategy is optimal for $\Delta t \to 0$ and $V_{max} \to 0$.

Note that the underlying assumption in Theorem 4 is that partitions $\bigcup_{q \in Q_{safe}} q = S_{safe}$. If $Q_{safe}$ is a conservative under approximation of $S_{safe}$, i.e., $\bigcup_{q \in Q_{safe}} q \subset S_{safe}$, the lower bound in (38) still holds, but the upper bound may be smaller than the actual upper bound. Nevertheless, both bounds always hold for $Q_{safe}$, i.e., $P_{safe}(s, Q_{safe}, \tau \mid \sigma^*_{\mathcal{H}}, s_0) \in \left[ \hat{p}^k_{\sigma^*}(q_0), \hat{p}^k_{\sigma^*}(q_0) \right]$. Also note that, for an initial continuous state $x(0) = x_0$, there are $\vert A \vert$ choices for the initialization of the hybrid system $\mathcal{H}$, i.e., the choice of $a_0$ in $s_0 = (a_0, x_0)$. Let $q_0^a$ denote the region that contains $x_0$ in mode $a \in A$. Then, the optimal choice of mode for $s_0$ is the one that corresponds to the IMDP state with maximum lower bound, i.e., $a_0 = \arg \max_{a \in A} \hat{p}^k_{\sigma^*}(q_0^a)$.

An interesting question in the analysis of the proposed framework is how changes in the discretization variables affect the error. In the case of time discretization, Theorem 3 guarantees that the error due to the continuous transition kernel $T^c$ for a single transition step goes to zero exponentially with $\Delta t$. However, a smaller $\Delta t$ causes the distribution of the discrete transition kernel $T^d$ to have a smaller variance, which can lead to an increase in the space discretization error due to
Therefore, in order to obtain a smaller overall error, a finer space discretization is required, causing an increase in the number of states in the abstraction IMDP. This gives rise to a trade-off in error contribution by time and space discretization, which is empirically studied in the next section. A thorough analysis on how to tune the discretization parameters is the subject of future work.

Finally, it is important to note that, in order to obtain an IMDP abstraction of process \( x \), a hybrid system modeling of \( x \) is not necessarily required. One can always partition \( X_{\text{safe}} \) and compute the transition probability bounds \( \hat{P} \) and \( \tilde{P} \) through evaluations of \( T^c \) and \( T^d \) as in (20) and (21). The computations for \( T^c \) and \( T^d \), however, are difficult for arbitrary geometry of discrete cells. As a matter of fact, this step is known to be the bottleneck of verification and synthesis for continuous stochastic processes. In our approach, we overcome this burden by a particular choice of geometry for the discretization that is dependent on the dynamics of the stochastic process under each action. The hybrid system modeling allows this unique discretization of the stochastic process under each action. The hybrid system for the discretization that is dependent on the dynamics of the boundary, \( \partial X \), is a much smaller burden given the low (polynomial) computational complexity of the IMDP synthesis algorithm.

VII. Experimental Results

We implemented our abstraction and synthesis algorithms in MATLAB and tested the performance of the framework on a set of case studies. We first considered a stochastic process with a single mode to illustrate the efficacy and scalability of the approach in the formal analysis of SDEs (Case Study 1). Then, we analyzed the trade-off between space and time discretization parameters and their effects on the overall error (Case Study 2). Finally, we considered a strategy synthesis problem on a two-mode switched stochastic process (Case Study 3). The case studies presented in this section have been distilled as particular instances of a model for a multi-room heating system, where the continuous dynamics represent the evolution of the temperature in a set of rooms (with number equal to the state space dimension), and the discrete modes represent the different heating actions, which may be turned on or off according to the selected strategy [20]. All the computations were performed on an Intel Dual Core i5 machine with 8 GB of RAM using a single thread program. Nevertheless, the implementation of the abstraction construction is highly parallelizable, allowing significant speedups.

A. Case Study 1 - Verification

We consider a stochastic process in the form of (1) with \( A = \{a_1\} \),

\[
F(a_1) = \begin{pmatrix} -1 & 0 \\ 0 & -0.5 \end{pmatrix}, \quad G(a_1) = I,
\]

and \( X_{\text{safe}} = [-8, 8] \times [-8, 8] \). We are interested in the computation of the safety probability for every possible initial state of this model for the duration \( \tau = 1 \). In order to perform this analysis, we abstract the model to an IMDP. We use a sampling time of \( \Delta t = 0.1 \) and discretize \( X_{\text{safe}} \) in the transformed space induced by transformation function \( T_{a_1} \), i.e., \( \text{Post}(X_{\text{safe}}|T_{a_1}) \), by a grid such that the size of each cell in the original space is \( \Delta x = 0.5 \) per side. This results in an IMDP abstraction with \( |Q| = 1025 \) states, and the discrete safety time steps \( k_r = 10 \). We then run our IMDP (synthesis) algorithm to obtain the lower and upper bounds of the probability of safety, i.e., \( \hat{\rho}^{10}(q) \) and \( \tilde{\rho}^{10}(q) \), for each possible initial state \( q \in Q \).

Figure 2a shows \( X_{\text{safe}} \) and the space discretization, where the lower bound of the probability of safety \( \hat{\rho}^{10} \) is marked as a shade of gray. The obtained upper bound probabilities were \( \tilde{\rho}^{10}(q) \geq 0.97 \) for all \( q \in Q \). It took a total of 175 seconds to compute the abstraction and the safety probability bounds.

In order to thoroughly analyze the error introduced by the discretization variables, we consider the local error term

\[
\varepsilon_q = \hat{\rho}^{k_r}(q) - \tilde{\rho}^{k_r}(q),
\]

and introduce the error metrics of \( \varepsilon_{\text{med}} \), which is the median of \( \varepsilon_q \) for all \( q \in Q \), and \( \varepsilon_{\text{ave}} \), which is the average of \( \varepsilon_q \) per unit volume given by

\[
\varepsilon_{\text{ave}} = \frac{\sum_{q \in Q} V_q \varepsilon_q}{\sum_{q \in Q} V_q},
\]

where \( V_q \) is the area (volume) of the region associated to state \( q \). For the chosen discretization variables \( \Delta t = 0.1 \) and \( \Delta x = 0.5 \), we obtain \( \varepsilon_{\text{med}} = 0.08 \) and \( \varepsilon_{\text{ave}} = 0.35 \). The smaller median error (than the average error) means that at least half of the cells have small errors \( \varepsilon_q \leq 0.08 \), and a small number of the cells have very large errors. This can be explained by analyzing Figure 2a. It is easy to observe that \( \hat{\rho}^{10}(q) \approx 0 \) for the cells \( q \in Q \) near \( \partial X_{\text{safe}} \), the boundary of \( X_{\text{safe}} \), i.e., \( \varepsilon_q \approx 1 \). In turn, for the cells \( q \in Q \) that are not close to the boundary, \( \varepsilon_q \approx 0 \). This observation is aligned with the theory and the intuition as expressed in (32) and (33).

That is, \( x \) is a continuous-time diffusion process with a stable drift term \( F(a_1) \) (the real eigenvalues are negative); hence, the process gets a strong pull towards the center, resulting in large upper bounds for safety probability in snapshots of \( \Delta t \) (discrete transition kernel), as suggested by (32). However, during this time, the probability of its continuous trajectory leaving \( X_{\text{safe}} \) is high due to the stochastic nature of the process if \( x \) starts near \( \partial X_{\text{safe}} \), resulting in near-zero lower bound for the continuous transition kernel, i.e., \( \alpha \approx 0 \) in (33). This suggests that the use of smaller cells near \( \partial X_{\text{safe}} \) can potentially reduce the spread of this effect (large error) to the nearby states, i.e., adaptive (non-uniform) gridding, as in [11].

B. Case Study 2 - Analysis of Error Trade-off

In order to explore the trade-offs between the time and space discretization parameters and their effects on the overall error, we consider the same system as in Case Study 1 (Section VII-A) and perform safety analysis for the duration \( \tau = 1 \) with various values for \( \Delta t \in \{0.05, 0.10, 0.15, 0.20, 0.25\} \) and \( \Delta x \in \{0.20, 0.32, 0.40, 0.50, 0.64\} \). The obtained results are shown in Figure 3, illustrating the change in the distribution of \( \varepsilon_q \) for the different values of \( \Delta t \) and \( \Delta x \).
Figure 3a is the box plot of $\varepsilon_q$, showing the distribution of $\varepsilon_q \forall q \in Q$ for each pair $(\Delta x, \Delta t)$ along with $\varepsilon_{med}$ (indicated with a red line) and $\varepsilon_{ave}$ (indicated with a black star). A general observation is that, by decreasing $\Delta x$, the distribution of the error becomes narrower and its average value decreases. In Figure 3b, we plot $\varepsilon_{ave}$ as a function of $\Delta x$ and $\Delta t$. For each $\Delta t$, $\varepsilon_{ave}$ tends to decrease linearly with $\Delta x$. This is intuitive because a smaller $\Delta x$ means a smaller error due to space discretization for the same $\Delta t$. However, if we fix $\Delta x$ and consider different values of $\Delta t$, we have a substantially different behavior: the error curve becomes parabolic. That is, for each $\Delta x$, there is an optimal value of $\Delta t$ that minimizes $\varepsilon_{ave}$, and this value of $\Delta t$ is not necessarily the smallest one. This is due to the fact that, as discussed in Section VI, decreasing $\Delta t$ causes a smaller variance in the discrete transition kernel, which means smaller cells are required to obtain a smaller overall error. Moreover, the smaller the $\Delta x$, the smaller the optimal value of $\Delta t$ becomes. This means that, for a given finite $\Delta x$, if $\Delta t$ is chosen to be too small, some of the dynamics of the system may not be appropriately reflected in the abstraction, hence resulting in large errors.

This case study shows that the choice of values for discretization parameters $\Delta t$ and $\Delta x$ in order to obtain a desired overall error is not trivial. One possible approach in tuning these parameters is to adapt a refinement technique to iteratively refine $\Delta t$ and $\Delta x$ to arrive at the desired overall error value. The proposed framework enables such an approach due to the low computation cost of abstraction and synthesis algorithms. Another approach is to obtain a conservative approximation of the optimal value of $\Delta x$ for a given $\Delta t$ before performing the computations, as shown in [1]. However, since such approaches tend to be very conservative compared to our approach, this results in a very small $\Delta x$ to guarantee a given overall error, leading to state explosion.

C. Case Study 3 - Switching Strategy Synthesis

We add a second mode $a_2$ to the process in Case Study 1 in Section VII-A, i.e., $A = \{a_1, a_2\}$, where

$$F(a_2) = \begin{pmatrix} -0.5 & 0.1 \\ 0 & -1 \end{pmatrix}, \quad G(a_2) = I.$$ 

We are interested in synthesizing a switching strategy that maximizes the probability of safety at every possible initial state with the duration of $\tau = 1$ for the same $X_{safe}$ as in Case Study 1. In abstraction, we again use a sampling time $\Delta t = 0.1$. For space discretization, we use an adaptive grid such that the resulted cells have the maximum and minimum cells sizes of $\Delta x_{max} = 0.5$ and $\Delta x_{min} = 0.1$ in the original space. Our adaptive grid algorithm first over-approximates $Post(X_{safe}|T_{a_i})$ for $i \in \{1, 2\}$ by using a uniform grid with the allowed maximum-sized cells. Then, it refines the cells that are partly unsafe up to the resolution of the minimum-sized cells to under-approximate the safe set. Figures 2b and 2c show the discretization of modes $a_1$ and $a_2$, respectively. Note that, in mode $a_2$, the cells are rotated which indicate that the transformation function $T_{a_2}$ includes a rotation in addition to translation. The IMDP has a total of $|Q| = 2947$ states with $|Q^{a_1}| = 1024$, $|Q^{a_2}| = 1922$, and one $q_{unsafe}$. The total abstraction took 54 minutes.

We then run our synthesis algorithm to obtain the robust strategy $\sigma^*$ with its corresponding safety probability bounds $\hat{p}^{\sigma^*}_q(q)$ and $\check{p}^{\sigma^*}_q(q)$ for all $q \in Q$. In Figures 2b and 2c the lower bounds $\check{p}^{\sigma^*}_q(q)$ are shown as shades of gray. The upper bounds were $\hat{p}^{\sigma^*}_q(q) \geq 0.93 \forall q \in Q$. The synthesis computation took 15 seconds.

Call $\bar{\sigma}$ the strategy that always picks action $a_1$. This strategy gives rise to the results obtained in Case Study 1 and presented in Figure 2a. Comparing it with the results of the strategy $\sigma^*$ in Figure 2b, it becomes evident that $\sigma^*$ improves the safety probability bounds, especially for the states near the top and bottom sides of $X_{safe}$. This means that $\sigma^*$ takes advantage of the drift in mode $a_2$ to keep the process safe in those sections of $X_{safe}$. More precisely, the median and average errors obtained for mode $a_1$ under $\sigma^*$ are $\varepsilon_{med} = 0.01$ and $\varepsilon_{ave} = 0.28$, which show a reduction from the errors under $\bar{\sigma}$ where $\varepsilon_{med} = 0.08$ and $\varepsilon_{ave} = 0.35$.

VIII. CONCLUSIONS

In this work, we proposed a theoretical and computational framework for analysis and synthesis for switched stochastic systems that evolve in continuous time. The framework mitigates the problem of state explosion through a suitable choice
of the abstraction model and the derivation of discretization methods for both time and space domains that result in tight error bounds (exact for space). The latter is specifically enabled by stochastic hybrid system modeling and a novel dynamic-dependent space discretization. This leads to fine and compact abstractions, whose computation is fast. Even though the framework is presented for synthesis problems with safety properties, it can be extended, in a straightforward manner, to verification and synthesis for more complex and even multi-objective properties expressed in, e.g., PCTL and CSL.

One of the main results of this study is the inherent trade-off in the error contribution by the time and space discretization parameters. An empirical analysis of this trade-off is performed in this work. An interesting direction for future work is a thorough analysis on how the discretization parameters can be tuned to obtain a desired overall error. One potential candidate is sequential adaptive gridding in space discretization and iterative refinement of the parameters.

REFERENCES


However, since the set of vertices is finite, it is generally faster to just include all the vertices as possible candidate solutions instead of solving the system of equations.

**Case 3:** \( y \) is in the boundary of \( \text{Post}(q'_i|T_a) \), but is not a vertex. In this case only \( r < m \) of the half-spaces in \( H \) intersect at \( y \). Thus, if \( y \) is a maximum then \( \nabla f(y) = H^T \mu \), where \( H \) is the submatrix of \( H \) containing the \( r < m \) half-spaces intersecting at \( y \), and \( \mu \) contains only the \( r \) corresponding constants. Note that this is a system with more equations than variables. Therefore, only when some of the constraints become linearly dependent, there may be a solution for \( y \) in \( \text{Post}(q'_i|T_a) \), if at all.

The minimum case is identical except that condition \( \nabla f(y) = H^T \mu \) is replaced with \( \nabla f(y) = -H^T \mu \).

**APPENDIX B**

**Proof of Theorem 3.** By definition of \( T^c(\Delta|q_j, x_1, a) \) we have that

\[
\min_{x_1 \in q_j} T^c(\Delta|q_j, x_1, a) \leq 1 - \max_{(x_1, x_2) \in (q_1, q_2)} \text{Prob}^b_x(\exists t \in [0, \Delta t], b_x(t) \not\in X_{safe}) \\
\geq 1 - \max_{(x_1, x_2) \in (q_1, q_2)} \text{Prob}^b_x(\exists t \in [0, \Delta t] \text{ s.t. } |b_x(t) - x_1| \geq \epsilon_q, \text{ and } |b_x(t) - x_2| \geq \epsilon_q). \]

The last inequality exploits the fact that each path of \( b_x \) that reaches a state outside \( X_{safe} \) during \([0, \Delta t]\) travels a distance of at least \( \epsilon_q = \max\{\epsilon_{q_1}, \epsilon_{q_2}\} \). Without any loss of generality, assume \( \epsilon_q = \epsilon_q \). Then, we have

\[
\min_{x_1 \in q_j} T^c(\Delta|q_j, x_1, a) \geq 1 - \max_{(x_1, x_2) \in (q_1, q_2)} \text{Prob}^b_x(\sup_{t \in [0, \Delta t]} |b_x(t) - x_1| \geq \epsilon_q) \\
\geq 1 - \max_{(x_1, x_2) \in (q_1, q_2)} \sum_{i \in [1, \ldots, m]} \text{Prob}^b_x(\sup_{t \in [0, \Delta t]} |b_x^{(i)}(t) - x_1^{(i)}| \geq \epsilon_{q/m}).
\]

The last inequality holds due to the union bound. By the linearity of Gaussian processes, we can write \( b_x(t) = E_{b_x}(t) + b_{x}(t) \), where \( b_x(t) \) is a normally distributed random variable with zero-mean and covariance equal to the one of \( b_x(t) \). By the application of the triangle inequality, we then have

\[
\min_{x_1 \in q_j} T^c(\Delta|q_j, x_1, a) \geq 1 - \max_{(x_1, x_2) \in (q_1, q_2)} \sum_{i \in [1, \ldots, m]} \text{Prob}^b_x(\sup_{t \in [0, \Delta t]} |E_{b_x}^{(i)}(t) - x_1^{(i)}| + |b_{x}^{(i)}(t_1)| \geq \epsilon_{q/m} ) \\
= 1 - \max_{(x_1, x_2) \in (q_1, q_2)} \sum_{i \in [1, \ldots, m]} \text{Prob}^b_x(\sup_{t \in [0, \Delta t]} |E_{b_x}^{(i)}(t_1)| \geq m \epsilon_{q/m}) \geq 1 - 2 \sum_{i \in [1, \ldots, m]} \text{Prob}^b_x(\sup_{t \in [0, \Delta t]} |E_{b_x}^{(i)}(t_1)| \geq \epsilon_{q,i}).
\]

**APPENDIX A**

**Proof of Theorem 1.** We first consider the maximum case and then discuss the minimum case. The KKT conditions guarantee that if \( y \in \text{Post}(q'_i|T_a) \) is a local maximum for \( f \), then there must exist a vector of constants \( \mu = (\mu_1, \ldots, \mu_k) \) such that \( \nabla f(y) = H^T \mu \), \( \mu_i \geq 0 \) for all \( i \in \{1, \ldots, k\} \), and \( \mu_i(H(i,j)y(j) - b_i) = 0 \), where \( H(i,j) \) is the component in the \( i \)-th row and \( j \)-th column of matrix \( H \). Note that we have a constant \( \mu_i, i \in \{1, \ldots, k\} \), for each of the half-spaces defining \( \text{Post}(q'_i|T_a) \). Thus, there are three possible cases:

**Case 1:** \( x^* \) is not in the boundary of \( \text{Post}(q'_i|T_a) \). In this case the KKT conditions imply that \( y \) is a maximum only if \( \nabla f(y) = 0 \). For a normal distribution with identity covariance, this point is exactly \( y = (\frac{1}{2}v^{(1)}+\frac{1}{2}v^{(2)}, \ldots, \frac{1}{2}v^{(n)}+\frac{1}{2}v^{(m)}) \). If \( y \in \text{Post}(q'_i|T_a) \), then this is the global maximum, because it is the global maximum of the unconstrained problem.

**Case 2:** \( x^* \) is in the boundary of \( \text{Post}(q'_i|T_a) \). We call a vertex an intersection of \( m \) half-spaces. As a consequence, we have that the KKT conditions are satisfied in \( y \), vertex of \( \text{Post}(q'_i|T_a) \), if and only if \( \nabla f(y) = H^T \mu \), where \( H \) is the submatrix that contains only the \( m \) rows of \( H \) representing the half-spaces interesting at \( y \), and \( \mu \) contains only the \( m \) corresponding constants. Thus, we have a system of \( m \) equations and \( m \) variables that has solution for \( \mu_i \in \mathbb{R} \).
where $\bar{e}_{q,i} = \frac{e_q}{m} - L_{b(i)}$ for $L_{b(i)} = \sup_{(x_1,x_2) \in [q_1,q_2]} |E_{b(i)}^x(t) - x_i^i(t)|$. The last inequality holds because $Cov_{b(i)}(t,s)$ is independent of $x_1,x_2$. Thus, the only term depending on $x_1,x_2$ is $e_q$. Since $b(i)^x(t)$ is a uni-dimensional Gaussian Process (GP), we can bound its supremum during $[0,\Delta t]$ by using the Borell-TIS inequality [37]. For $D > E \left( \sup_{t \in [0,\Delta t]} b(i)^x(t) \right)$, the Borell-TIS inequality [37] guarantees that

\[
Prob_{\bar{b}} \left( \sup_{t \in [0,\Delta t]} b(i)^x(t) > D \right) \leq e^{-\frac{D - E \left( \sup_{t \in [0,\Delta t]} b(i)^x(t) \right)^2}{2e^{\left(\bar{e}_{q,i} - Q(i)^{x} \right)^2}}},
\]

where $e^{\left(\bar{e}_{q,i} - Q(i)^{x} \right)} = \sup_{t \in [0,\Delta t]} Cov_{b(i)}(t)$.

In order to bound $E \left( \sup_{t \in [0,\Delta t]} b(i)^x(t) \right)$, we consider the pseudometric $d_i$ and constant $K_{d,i}^2$ as defined in Sec. IV-D. By the Dudley’s entropy integral [37], we obtain

\[
E \left( \sup_{t \in [0,\Delta t]} b(i)^x(t) \right) \leq 12 \int_0^{K_{d,i}^2 \Delta t} \ln \left( \frac{2K_{d,i}^2 \Delta t}{x} + 1 \right) dx.
\]

By substituting this expression in the Borell-TIS inequality and picking $D = \frac{e_q}{m}$, we arrive to the result.

**APPENDIX C**

**Proof of Proposition 3.** For the upper bound, we have that for $q_i \in Q_{safe}$ and $a \in A$,

\[
\max_{x \in q_i} P_{safe}(X_{safe}^{\Delta t} | x, a) 
\leq \left( \max_{x \in q_i} T^d(X_{safe} | x, a, \Delta t) \right) \left( \max_{x \in q_i} T^c(\Delta t | x, X_{safe}) \right)
\leq \max_{x \in q_i} T^d(X_{safe} | x, a, \Delta t)
\leq \max_{x \in q_i} \int_{X_{safe}} N(z | E_x^{\Delta t}, Cov_x(\Delta t)) dz
\leq \max_{y \in Post(q_i | T_a)} \sum_{q \in Q^n} \int_{Post(q | T_a)} N(z | y, 1) dz
= \max_{y \in Post(q_i | T_a)} \sum_{q \in Q^n} f(y, q).
\]

For the lower bound, we have that

\[
\min_{x \in q_i} P_{safe}(X_{safe}^{\Delta t} | x, a) 
\geq \left( \min_{x \in q_i} T^d(X_{safe} | x, a, \Delta t) \right) \left( \min_{x \in q_i} T^c(\Delta t | x, Q_{safe}, a) \right).
\]

Similarly to the upper bound, we have that

\[
\min_{x \in q_i} T^d(X_{safe} | x, a, \Delta t) \geq \min_{y \in Post(q_i | T_a)} \sum_{q \in Q^n} f(y, q).
\]

For $T^c$, by using a similar reasoning as in the one in Theorem 3, we have that

\[
\min_{x \in q_i} T^c(\Delta t | x, Q_{safe}, a) \geq \begin{cases} 
\max \{0, 1 - 2 \sum_{i=1}^m e^{-\frac{e_q}{2\Delta x}} \} & \text{if } \eta_{x,i} > 0, \forall i \in \{1, \ldots, m\} \\
0 & \text{otherwise}
\end{cases}
\]

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